

# The classification of symmetric conservative clones with a finite carrier and its applications in Computational Social Choice

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# Basic definitions

- Let  $A$  be a non-empty finite set.  $\mathcal{O}(A)$  is the set of all function (of any arity) on  $A$ . A set  $\mathcal{F} \subseteq \mathcal{O}(A)$  is a *clone* if it is closed under composition and contains all projections. The clone of all projections on  $A$  is denoted by  $\mathcal{E}(A)$ . The set of all  $n$ -ary function  $f \in \mathcal{F} \subseteq \mathcal{O}(A)$  is denoted by  $\mathcal{F}_{[n]}$ .
- A clone  $\mathcal{F} \subseteq \mathcal{O}(A)$  is *symmetric* if

$$f \in \mathcal{F} \rightarrow f_{\sigma} \in \mathcal{F}$$

for all function  $f \in \mathcal{O}(A)$  and permutation  $\sigma$  of  $A$ , where for any function  $f \in \mathcal{O}(A)_{[n]}$  the function  $f_{\sigma}$  is the  $n$ -ary function in  $\mathcal{O}(A)$  defined by

$$(\forall \mathbf{a} \in A^n) f_{\sigma}(\mathbf{a}) = \sigma^{-1} f(\sigma \mathbf{a}).$$

# Basic definitions

- A clone  $\mathcal{F} \subseteq \mathcal{O}(A)$  is *conservative* if it contains only conservative functions, i.e. functions  $f \in \mathcal{O}(A)$  satisfying

$$(\forall \mathbf{a} \in \text{dom} f) f(\mathbf{a}) \in \text{ran } \mathbf{a},$$

where  $\text{ran}(a_1, \dots, a_n) = \{a_1, \dots, a_n\}$  for any  $n < \omega$ .

## Note

*In terms of Galois connection  $(\text{Inv}, \text{Pol})$ , a conservative function is a function preserving any unary predicate. A Boolean function is conservative if it preserves  $\mathbf{0}$  and  $\mathbf{1}$ , i.e., belongs to Post's class  $T_{01}$ .*

# Objectives of the study and references

- An explicit classification of symmetric conservative clones on a finite set  $A$ .
- Description of their invariant sets  $\mathfrak{D} \subseteq A^Q$ .
- Application to Computational Social Choice.

## References

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- [2] Ježek J. and Quackenbush R., Minimal clones of conservative functions. Int. J. of Alg. and Comp., vol. 05, No. 06 (1995), pp. 615-630
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# The classification of symmetric conservative clones on a finite set $A$

Let  $\mathcal{F} \subseteq \mathcal{O}(A)$  be a clone. We define the characteristic  $\chi(\mathcal{F})$  that uniquely identifies  $\mathcal{F}$ .

- $r(\mathcal{F})$  is the minimal positive arity of a function  $f \in \mathcal{F} \setminus \mathcal{E}(A)$ .  
If  $\mathcal{F} = \mathcal{E}(A)$ , then  $r(\mathcal{F}) = \omega$ .
- $m(\mathcal{F})$  is a minimal positive integer  $m$  for which there are different  $a, b \in A$ , pairwise different  $c_1, \dots, c_m \in A$  and a function  $f \in \mathcal{F}$  such that

$$f(a, b, \dots, b) = a \text{ and } f(a, \dots, a, c_1, \dots, c_m) \neq a.$$

If the set of such  $m$  is empty, then  $m(\mathcal{F}) = \omega$ .

# The classification of symmetric conservative clones on a finite set $A$

- $R_n(\mathcal{F})$  is a binary relation on  $A^n$  defined by

$$\mathbf{a} R \mathbf{b} \leftrightarrow (\exists \sigma \in S_A)(\forall f \in \mathcal{F}_{[n]}) f(\mathbf{b}) = \sigma f(\mathbf{a}),$$

where  $S_A$  is the set of all permutations of  $A$ .

- $R(\mathcal{F}) = \bigcup_{n < \omega} R_n(\mathcal{F})$

A binary relation  $R$  on  $A^{<\omega}$  is called *stable* if for all  $\mathbf{a}, \mathbf{b} \in A^{<\omega}$

- $\mathbf{a} R \mathbf{b} \rightarrow \mathbf{b} = \sigma \mathbf{a}$  for some permutation  $\sigma$  of  $A$ ,
- $\mathbf{a} R \mathbf{b} \rightarrow \sigma \mathbf{a} \tau R \sigma \mathbf{b} \tau$  for any permutation  $\sigma$  of  $A$  and any function  $\tau : \{1, \dots, k\} \rightarrow \{1, \dots, n\}$ , where  $\mathbf{a}, \mathbf{b} \in A^n$

## Proposition

For any symmetric clone  $\mathcal{F} \subseteq \mathcal{O}(A)$  the relation  $R(\mathcal{F})$  is a stable equivalence relation on  $A^{<\omega}$ .

# The classification of symmetric conservative clones on a finite set $A$

For any clone  $\mathcal{F} \subseteq \mathcal{O}(A)$  and a set  $B \subseteq A$  the set  $\{f \upharpoonright B^{<\omega} : f \in \mathcal{F}\}$  is called the restriction of  $\mathcal{F}$  to  $B$ . A restriction of  $\mathcal{F}$  to an arbitrary  $B \subseteq A$  is a clone for any conservative clone  $\mathcal{F}$ .

## Proposition

*For any symmetric conservative clone  $\mathcal{F} \subseteq \mathcal{O}(A)$  there is exactly one Post's class  $P \subseteq T_{01}$  closed with respect to duality which is naturally equivalent to the restriction of  $\mathcal{F}$  to an arbitrary two-element set  $B \subseteq A$ .*

This Post's class is denoted by  $\Pi(\mathcal{F})$ .

- The quadruple  $\chi(\mathcal{F}) = (r(\mathcal{F}), m(\mathcal{F}), R(\mathcal{F}), \Pi(\mathcal{F}))$  is called the *characteristic* of a symmetric conservative clone  $\mathcal{F}$ .



# The classification of symmetric conservative clones on a finite set $A$

## Theorem

Let  $A$  be a finite set of cardinality  $|A| \geq 2$ . Then for any  $r, m < \omega + 1$ , stable equivalence relation  $R$  on  $A^{<\omega}$  and Post's class  $P \subseteq T_{01}$  closed with respect to duality there is not more than one symmetric conservative clone  $\mathcal{F} \subseteq \mathcal{O}(A)$  with  $\chi(\mathcal{F}) = (r, m, R, P)$ .

## Proposition

- There are quadruples  $(r, m, R, P)$  for which there is no clone  $\mathcal{F}$  with  $\chi(\mathcal{F}) = (r, m, R, P)$ .
- If  $r(\mathcal{F}) \geq 4$ , all function  $f \in \mathcal{F}_n$  coincide with a projection on a set  $A_{<r}^n = \{\mathbf{a} \in A^n : |\text{ran } \mathbf{a}| < r\}$ .

## Invariant sets of symmetric conservative clone

Let  $A$  and  $Q$  be finite sets. A function  $f \in \mathcal{O}(A)_{[n]}$  *preserves* a set  $\mathfrak{D} \subseteq A^Q$  if for all  $\mathfrak{d}_1, \dots, \mathfrak{d}_n \in \mathfrak{D}$  the set  $\mathfrak{D}$  contains the function  $f(\mathfrak{d}_1, \dots, \mathfrak{d}_n)$  defined by

$$f(\mathfrak{d}_1, \dots, \mathfrak{d}_n)(q) = f(\mathfrak{d}_1(q), \dots, \mathfrak{d}_n(q))$$

for all  $q \in Q$ .

For any set  $\mathfrak{D} \subseteq A^Q$ , function  $f \in \mathcal{O}(A)$ , and sets  $\mathbb{D} \subseteq \mathcal{P}(A^Q)$  and  $\mathcal{F} \subseteq \mathcal{O}(A)$

- $\text{Pol } \mathfrak{D} = \{f \in \mathcal{O}(A) : f \text{ preserves } \mathfrak{D}\};$
- $\text{Inv}_Q f = \{\mathfrak{D} \subseteq A^Q : f \text{ preserves } \mathfrak{D}\};$
- $\text{Pol } \mathbb{D} = \bigcap_{\mathfrak{D} \in \mathbb{D}} \text{Pol } \mathfrak{D};$
- $\text{Inv}_Q \mathcal{F} = \bigcap_{f \in \mathcal{F}} \text{Inv}_Q f.$

# Invariant sets of symmetric conservative clone

## Proposition

*The couple  $(\text{Inv}_Q, \text{Pol})$  is a Galois connection between the Boolean lattices  $\mathcal{P}(\mathcal{O}(A))$  and  $\mathcal{P}(\mathcal{P}(A^Q))$ . Galois-closed sets  $\mathcal{F} \subseteq \mathcal{O}(A)$  are clones.*

## Note

- If  $Q = \{1, \dots, m\}$ , then any set  $\mathcal{D} \subseteq A^Q$  is (identified with) the predicate  $\{(\mathfrak{d}(1), \dots, \mathfrak{d}(m)) : \mathfrak{d} \in \mathcal{D}\}$ .  
 $\mathcal{D} \in \text{Inv}_Q \mathcal{F}$  iff  $\mathcal{D}$  is an  $m$ -ary predicate in  $\text{Inv} \mathcal{F}$ .
- For any clone  $\mathcal{F} \subseteq \mathcal{O}(A)$ ,  $\mathcal{F}_{[n]} \in \text{Inv}_{A^n} \mathcal{F}$ .

# Invariant sets of symmetric conservative clone

Notation: for any  $\mathfrak{D} \in A^Q$  and  $q \in Q$ ,  $\mathfrak{D}(q) = \{\mathfrak{d}(q) : \mathfrak{d} \in \mathfrak{D}\}$ .

## Theorem

Let  $A$  and  $Q$  be non-empty finite sets, and  $|A| \geq 2$ . Let  $\mathcal{F}$  be a symmetric conservative clone on  $A$ . Then

1. if  $r(\mathcal{F}) \geq 4$  or  $r(\mathcal{F}) = 3 \wedge \Pi(\mathcal{F}) = O_1$  (where  $O_1$  is a class of all Boolean projections), then any set  $\mathfrak{D} \in \text{Inv}_Q \mathcal{F}$  is an intersection of
  - a set of the form  $\{\mathfrak{d} \in A^Q : \mathfrak{d} \upharpoonright P \in \mathfrak{C}\}$  where  $P \subseteq Q$ , and  $\mathfrak{C}$  is a subset of  $A^P$  satisfying  $(\forall q \in P) |\mathfrak{C}(q)| < r(\mathcal{F})$ ,
  - a finite number of sets of the form  $\{\mathfrak{d} \in A^Q : \mathfrak{d}(q) \in B\}$  where  $q \in Q$  and  $B \subseteq A$ , and
  - a finite number of sets of the form  $\{\mathfrak{d} \in A^Q : \mathfrak{d}(q) = \sigma \mathfrak{d}(p)\}$  where  $p, q \in Q$  and  $\sigma \in S_A$ .

# Invariant sets of symmetric conservative clone

## Theorem (continued)

2. if  $r(\mathcal{F}) = 3$  and  $\Pi(\mathcal{F}) \notin \{O_1, L_4\}$  (where  $L_4$  is a class of all linear self-dual functions in  $T_{01}$ ), then any set  $\mathfrak{D} \in \text{Inv}_Q \mathcal{F}$  is an intersection of
- a finite number of sets of the form  $\{\mathfrak{d} \in A^Q : \mathfrak{d}(q) \in B\}$  where  $q \in Q$  and  $B \subseteq A$ , and
  - a finite number of sets of the form  $\{\mathfrak{d} \in A^Q : \mathfrak{d}(q) = \sigma \mathfrak{d}(p)\}$  where  $p, q \in Q$  and  $\sigma \in S_A$ , and
  - a finite number of sets of the form  $\{\mathfrak{d} \in A^Q : \mathfrak{d}(p) = b \vee \mathfrak{d}(q) = c\}$  where  $p, q \in Q$  and  $a, b \in A$ ;

## Theorem (continued)

3. if  $r(\mathcal{F}) = 3$  and  $\Pi(\mathcal{F}) = L_4$ , then any set  $\mathfrak{D} \in \text{Inv}_Q \mathcal{F}$  is an intersection of
- a finite number of sets of the form  $\{\mathfrak{d} \in A^Q : \mathfrak{d}(q) \in B\}$  where  $q \in Q$  and  $B \subseteq A$ , and
  - a finite number of sets of the form  $\{\mathfrak{d} \in A^Q : \mathfrak{d}(q) = \sigma \mathfrak{d}(p)\}$  where  $p, q \in Q$ ,  $\sigma \in S_A$ , and
  - a set  $\mathfrak{C} \in \text{Inv}_Q \mathcal{F}_0$  satisfying  $(\forall q \in Q) |\mathfrak{C}(q)| = 2$ , where  $\chi(\mathcal{F}_0) = (3, 1, L_4, R_0)$ ,  
 $R_0 = \text{Eq} \cup \{(\mathbf{a}, \mathbf{b}) : |\text{ran } \mathbf{a}| = |\text{ran } \mathbf{b}| \leq 2 \wedge (\exists \sigma \in S_A) \mathbf{b} = \sigma \mathbf{a}\}$ ,

## Theorem (continued)

4. if  $r(\mathcal{F}) = 2$ , then any set  $\mathfrak{D} \in \text{Inv}_Q \mathcal{F}$  is an intersection of
- a finite number of sets of the form  $\{\mathfrak{d} \in A^Q : \mathfrak{d}(q) \in B\}$  where  $q \in Q$  and  $B \subseteq A$ , and
  - a finite number of sets of the form  $\{\mathfrak{d} \in A^Q : \mathfrak{d}(q) = \mathfrak{d}(p)\}$  where  $p, q \in Q$ ,  $\sigma \in S_A$ , and
  - a finite number of sets of the form  $\{\mathfrak{d} \in A^Q : \mathfrak{d} \upharpoonright P \in \text{Inv}_P \mathcal{F}_B\}$  where  $P \subseteq Q$ ,  $B \subseteq A$ ,  $|B| = 2$ , and  $\mathcal{F}_B$  is the restriction of  $\mathcal{F}$  to  $B$ .

## Note

$r(\mathcal{F}) \geq 2$  for any conservative clone  $\mathcal{F}$ .

# Application to Computational Social Choice

- $A$  = a non-empty finite set (of alternatives);
- $k$  = a natural number (technical parameter),  $k \geq 1$ ;
- $[A]^k = \{B \subseteq A : |B| = k\}$ ;
- $k$ -choice function on  $A$  is a function  $\mathfrak{c} : [A]^k \rightarrow A$  satisfying

$$(\forall p \in [A]^k) \mathfrak{c}(p) \in p.$$

$k$ -choice function = *individual preferences*.

- $\mathfrak{C}_k(A)$  – the set of all  $k$ -choice functions on  $A$ ;
- $k$ -choice function  $\mathfrak{c}$  is rational if  $\mathfrak{c}(p) = \max_{\prec} p$  for some linear order on  $A$ ;
- $\mathfrak{R}_k(A) =$  the set of all rational  $k$ -choice functions on  $A$ ;



# Application to Computational Social Choice

A set  $\mathfrak{D} \subseteq \mathfrak{C}_k(A)$  is *symmetric* if for any function  $\mathfrak{c} \in \mathfrak{D}$  and permutation  $\sigma \in S_A$  the function  $\mathfrak{c}_\sigma$  defined by

$$(\forall p \in [A]^k) \mathfrak{c}_\sigma(p) = \sigma^{-1} \mathfrak{c}(\sigma p),$$

belongs to  $\mathfrak{D}$ .

- $n$  is a natural number (of voters),  $n \geq 1$ ;
- *social welfare function* (SWF) is a function

$$\varphi : (\mathfrak{C}_k(A))^n \rightarrow \mathfrak{C}_k(A)$$

# Application to Computational Social Choice

An  $n$ -ary SWF  $\varphi$  satisfies the *Arrow's conditions* iff

1. For all  $\mathbf{c}_1, \dots, \mathbf{c}_n \in \mathfrak{C}_k(A)$  and  $q \in [A]^k$

$$\varphi(\mathbf{c}_1, \dots, \mathbf{c}_n)(q) \in \{\mathbf{c}_1(q), \dots, \mathbf{c}_n(q)\}$$

2. For all  $\mathbf{c}_1, \dots, \mathbf{c}_n, \mathbf{c}'_1, \dots, \mathbf{c}'_n \in \mathfrak{C}_k(A)$  and  $q \in [A]^k$ .

$$\begin{aligned} (\mathbf{c}_1(q), \dots, \mathbf{c}_n(q)) = (\mathbf{c}'_1(q), \dots, \mathbf{c}'_n(q)) \rightarrow \\ \varphi(\mathbf{c}_1, \dots, \mathbf{c}_n)(q) = \varphi(\mathbf{c}'_1, \dots, \mathbf{c}'_n)(q) \end{aligned}$$

3. For all  $\mathbf{c}_1, \dots, \mathbf{c}_n \in \mathfrak{C}_k(A)$  and  $p, q \in [A]^k$

$$\begin{aligned} (\mathbf{c}_1(p), \dots, \mathbf{c}_n(p)) = (\mathbf{c}_1(q), \dots, \mathbf{c}_n(q)) \rightarrow \\ \varphi(\mathbf{c}_1, \dots, \mathbf{c}_n)(p) = \varphi(\mathbf{c}_1, \dots, \mathbf{c}_n)(q) \end{aligned}$$

# Application to Computational Social Choice

## Proposition

An  $n$ -ary SWF  $\varphi$  satisfies the Arrow's conditions iff there is a conservative function  $f : A^n \rightarrow A$  such that

$$\varphi(\mathbf{c}_1, \dots, \mathbf{c}_n)(p) = f(\mathbf{c}_1(p), \dots, \mathbf{c}_n(p)).$$

for all  $p \in [A]^k$  and  $\mathbf{c}_1, \dots, \mathbf{c}_n \in \mathfrak{C}_k(A)$ .

A set  $\mathfrak{D} \subseteq \mathfrak{C}_r(A)$  *has the Arrow property* if any SWF  $\varphi \in \text{Pol } \mathfrak{D}$  satisfying the Arrow's conditions is a projection. Equivalently, a set  $\mathfrak{D} \subseteq \mathfrak{C}_r(A)$  has the Arrow property if for any positive integer  $n$  any  $n$ -ary function  $f \in \text{Pol } \mathfrak{D}$  coincides with a projection on a set  $A_{\leq k}^n = \{\mathbf{a} \in A^n : |\text{ran } \mathbf{a}| \leq k\}$ .

# Application to Computational Social Choice

Theorem (K. Arrow, 1950, 1963)

*Let  $A$  be a finite set,  $|A| \geq 3$ . Then  $\mathfrak{R}_2(A)$  has the Arrow property.*

Theorem (S. Shelah, 2005)

*Let  $A$  be a finite set. Then there are natural numbers  $k_1, k_2$  (e.g.  $k_1 = k_2 = 7$ ) such that for any natural number  $k$ ,  $k_1 \leq r \leq |A| - k_2$ , any non-empty proper symmetric subset  $\mathfrak{D}$  of the set  $\mathfrak{C}_k(A)$  has the Arrow property.*

We give a complete description of symmetric sets  $\mathfrak{D} \subseteq \mathfrak{C}_k(A)$  with Arrow property.

# Application to Computational Social Choice

Let  $|A| = 4$  and let  $K$  be the *Klein four-group* of permutations of  $A$ . For any sets  $p, q \in [A]^3$  there is only one permutation  $\sigma_{p,q} \in K$  for which

$$q = \sigma_{p,q}(p).$$

- $\mathfrak{C}_3^K(A)$  is the set of all functions  $\mathfrak{c} \in \mathfrak{C}_3(A)$  such that

$$\mathfrak{c}(q) = \sigma_{p,q}\mathfrak{c}(p) \text{ for all } p, q \in [A]^3.$$

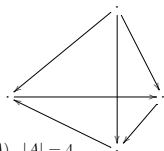
$q$	$\mathfrak{c}_0(q)$	$\mathfrak{c}_1(q)$	$\mathfrak{c}_2(q)$
$\{a, b, c\}$	$a$	$b$	$c$
$\{a, b, d\}$	$b$	$a$	$d$
$\{a, c, d\}$	$c$	$d$	$a$
$\{b, c, d\}$	$d$	$c$	$b$

# Application to Computational Social Choice

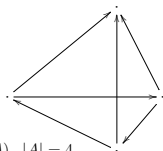
Let  $r = 2$  and  $|A| \geq 2$ . Any function  $\mathfrak{c} \in \mathfrak{C}_2(A)$  may be represented by the *tournament*  $\Gamma_{\mathfrak{c}} = (A, E)$  where

$$E = \{(a, b) \in A^2 : a \neq b \wedge \mathfrak{c}(\{a, b\}) = b\}.$$

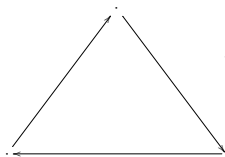
- The sets  $\mathfrak{C}_2^0(A)$  and  $\mathfrak{C}_2^1(A)$  are the sets of all functions  $\mathfrak{c} \in \mathfrak{C}_2(A)$  such that the *indegree* of any node of the tournament  $\Gamma_{\mathfrak{c}}$  is even (respectively, odd).



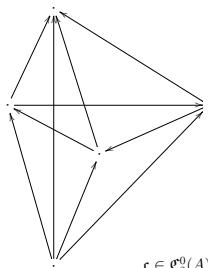
$c \in \mathfrak{C}_2^0(A), |A| = 4$



$c \in \mathfrak{C}_2^1(A), |A| = 4$



$c \in \mathfrak{C}_2^1(A), |A| = 3$



$c \in \mathfrak{C}_2^0(A), |A| = 5$

# Application to Computational Social Choice

## Theorem (P., 2014)

Let  $A$  be a finite set,  $r$  a natural number, and  $\mathcal{D}$  a non-empty proper symmetric subset of the set  $\mathfrak{C}_r(A)$ . Then the set  $\mathcal{D}$  does not have the Arrow property if and only if one of the following conditions holds:

- 1  $r = 2$ ,  $|A|$  equals 0 or 1 (mod 4), and  $\mathcal{D} = \mathfrak{C}_2^0(A)$ ,
- 2  $r = 2$ ,  $|A|$  equals 0 or 3 (mod 4), and  $\mathcal{D} = \mathfrak{C}_2^1(A)$ ,
- 3  $r = 2$ ,  $|A| = 0$  (mod 4), and  $\mathcal{D} = \mathfrak{C}_2^0(A) \cup \mathfrak{C}_2^1(A)$ ,
- 4  $r = 3$ ,  $|A| = 4$ , and  $\mathcal{D} = \mathfrak{C}_3^K(A)$ .

## Note (P., 2017)

The theorem remains true without condition 3 of Arrow's conditions.



Thank you!