

FUZZY SETS ON S AS CLOSURE OPERATIONS ON $P(S)$

B. Šešelja, G. Vojvodić

Prirodno-matematički fakultet. Institut za matematiku
21000 Novi Sad, ul.dr Ilije Djuričića br.4, Jugoslavija

ABSTRACT

It is known that a fuzzy set \bar{A} on S ($S \neq \emptyset$), as a mapping from S onto a complete lattice L , determines a family $\{A_p; p \in L\}$ of subsets of S indexed by the elements of L (i.e. $\bar{A} = \bigcup_{p \in L} p \cdot A_p$, see for example [1]).

In [2] it is proved that the lattice of all L -valued fuzzy sets on S is isomorphic with the lattice of specially constructed mappings from L to $P(S)$. Using a similar method, we prove that the partially ordered set $\bar{A}_L = \langle \{A_p; p \in L\}, \subseteq \rangle$ is a complete lattice - a quotient relative to one closure operation in $\langle P(S), \subseteq \rangle$. The converse is also proved, i.e. that a special complete lattice \bar{A}_L in $P(S)$, indexed by the elements of L , determines one fuzzy set \bar{A} on S . Moreover, we prove that every fuzzy set $\bar{A} : S \rightarrow L$ determines a complete lattice $L_{\bar{A}}$, as a quotient of a closure in L , and that $L_{\bar{A}}$ is isomorphic with a dual lattice of \bar{A}_L . This proves that every fuzzy set $\bar{A} : S \rightarrow L$ is, up to the lattice isomorphism, equal to one fuzzy set $\check{A} : S \rightarrow P(S)$.

We apply these results to the construction of a fuzzy

congruence relation $\bar{\theta}$ on an algebra A ([3]), using the lattice of all congruence relations on A .

0. Let $L = \langle L, \wedge, \vee, 0, 1 \rangle$ be a complete lattice, and $\bar{A} : S \rightarrow L$ a fuzzy set on S . Then

0.1. ($|1|$) $\bar{A} = \bigcup_{p \in L} p \cdot A_p$. Here we use the following:

a) For every $p \in L$, $A_p \subseteq S$, and $x \in A_p$ iff $\bar{A}(x) \geq p$ ($x \in S$);

b) We identify A_p with its characteristic function $A_p : S \rightarrow L$: if $x \in S$, then

$$A_p(x) = \begin{cases} 1, & \text{if } x \in A_p; \\ 0, & \text{otherwise;} \end{cases}$$

c) If $p \in L$, $x \in S$, then

$$p \cdot A_p(x) \stackrel{\text{def}}{=} p \wedge A_p(x).$$

d) Using c), we get the mapping

$$\bar{A}_p : S \rightarrow L, \bar{A}_p = p \cdot A_p, \text{ i.e. for } x \in S,$$

$$\bar{A}_p(x) = \begin{cases} p, & \text{if } x \in A_p, \\ 0, & \text{otherwise;} \end{cases}$$

e) The union of the family of fuzzy sets defined in d), i.e. $\bigcup_{p \in L} \bar{A}_p$ is the usual fuzzy union ([1]), that is, it is a mapping from S to L , such that for $x \in S$

$$\left(\bigcup_{p \in L} \bar{A}_p \right)(x) = \bigvee_{p \in L} \bar{A}_p(x).$$

Thereby, if $x \in S$

$$\bar{A}(x) = \bigvee_{p \in L} p \cdot A_p(x).$$

1. Let $\bar{A} : S \rightarrow L$ be an arbitrary fuzzy set on S. Let, also, $\bar{A}_L = \langle \{A_p; p \in L\}, \subseteq \rangle$ be a partially ordered set of subsets of S, determined by \bar{A} , as in 0.1. Then

1.1. The partially ordered set \bar{A}_L has the following properties:

- (1) $\bigcap_{p \in L_1} A_p = A_{\bigvee_{p \in L_1} p}$ (where L_1 is an arbitrary subset of L);
- (2) \bar{A}_L is a complete lattice in which the infimum is the intersection.

P r o o f.

(1) Let $x \in S$. Then, $x \in \bigcap_{p \in L_1} A_p$ iff for every $p \in L_1$ $\bar{A}(x) \geq p$ iff $\bar{A}(x) \geq \bigvee_{p \in L_1} p$ iff $x \in A_{\bigvee_{p \in L_1} p}$, since L is complete.

(2) Using (1), we get

(1') If for $p, q \in L$, $q \in L$, $p \leq q$, then $A_q \subseteq A_p$, since $p \leq q$ implies by (1) $A_p \cap A_q = A_{p \vee q} = A_q$. Now, from (1') it follows that for every $p \in L$,

$$0 \leq p \text{ implies } A_p \subseteq A_0, \text{ i.e.}$$

$A_0 = S$ is the greatest element in \bar{A}_L . Hence, using (1), we get that \bar{A}_L is a complete lattice.

The following statement is a direct consequence of the preceding one.

1.2. Define the mapping $A \rightarrow \hat{A}$ on $P(S)$, such that $\hat{A} = \bigcap_{A \subseteq A_p} A_p$. This mapping is a closure operation on $P(S)$, and the lattice \bar{A}_L is a quotient relative to that closure in a Boolean lattice $\langle P(S), \subseteq \rangle$.

P r o o f. A straightforward application of 1.1 and the wellknown properties of closure operations on lattices (see [5]).

The converse of 1.1 is the following proposition.

1.3. Let $A_p : L \rightarrow P(S)$ (for $p \in L$, $A_p \subseteq S$) be a mapping satisfying the following conditions:

- (1) $\bar{A}_L = \langle \{A_p; p \in L\}, \subseteq \rangle$ is the complete lattice ;
- (2) $\bigcap_{p \in L_1} A_p = A_{\bigvee_{p \in L_1} p}$ (for any subset L_1 of L);
- (3) $A_0 = S$. (see also [2]).

Let $\bar{A} = \bigcup_{p \in L} p \cdot A_p$, where $p \cdot A_p$ is defined as in c), 0.1., and the union is defined as in e), 0.1.

Then \bar{A} is a fuzzy set on S , and for $q \in L$

$$x \in A_q \quad \text{iff} \quad \bar{A}(x) \geq q.$$

P r o o f. \bar{A} is clearly a fuzzy set on S , that is a mapping $S \rightarrow L$, since $p \cdot A_p$ is for every $p \in L$ a fuzzy set by construction, and the union of fuzzy sets is again a fuzzy set.

Now, let $x \in S$, $q \in L$, and $x \in A_q$. Then

$\bar{A}(x) = \bigvee_{p \in L} p \cdot A_p(x) = \bigvee_{p \in L} \bar{A}_p(x) \geq q$, since $\bar{A}_q(x) = q$. Conversely, let $\bar{A}(x) \geq q$, and for $x \in S$, let $S_x = \{p; x \in A_p\}$. S_x is not empty, since by (3) $0 \in S_x$. Thus, $x \in \bigcap_{p \in S_x} A_p$. But

$$x \in \bigcap_{p \in S_x} A_p \quad \text{iff} \quad x \in A_{\bigvee_{p \in S_x} p} \quad \text{iff} \quad x \in A_{\bigvee S_x}.$$

We also have

$$q \leq \bigvee_{p \in L} p \cdot A_p(x) = \bigvee_{p \in L} \bar{A}_p(x) = \bigvee S_x, \quad \text{i.e.} \quad q \leq \bigvee S_x.$$

Using (2), we get $A_{\bigvee S_x} \subseteq A_q$, and since $x \in A_{\bigvee S_x}$, it

follows that $x \in A_q$.

Proposition 1.3 has following two simple corollaries.

1.4. If $A_p = \emptyset$, and $p \leq q$, then $A_q = \emptyset$.

1.5. If $A_q \subseteq A_p$, then $A_q = A_q \vee p$.

2. Let $\bar{A} = \bigcup_{p \in L} p \cdot A_p$ be an arbitrary fuzzy set on S. Define a relation \sim on L:

$$p \sim q \text{ iff } A_p = A_q.$$

2.1. \sim is an equivalence relation on L.

P r o o f. Straightforward.

2.2. Let $p, q \in L$. Then,

$$p \leq q \text{ implies } q \vee p \sim p.$$

P r o o f. By 1.5.

Define now a mapping $p \rightarrow \bigvee_{q \in |p|_{\sim}} q$ on L, and let $\bigvee_{q \in |p|_{\sim}} q = p_m$.

2.3. $p_m \in |p|_{\sim}$.

P r o o f. If $q \in |p|_{\sim}$, then

$$A_p = A_q = \bigcap_{q \in |p|_{\sim}} A_q = A \bigvee_{q \in |p|_{\sim}} q = A_{p_m}, \text{ i.e. } p \sim p_m.$$

Note that 2.3 is equivalent with $|p_m|_{\sim} = |p|_{\sim}$.

2.4. The mapping $p \rightarrow p_m$ is a closure operation on L.

P r o o f. 1° Since $p \in |p|_{\sim}$, it follows that

$$p \leq \bigvee_{q \in |p|_{\sim}} q = p_m, \text{ i.e. } p \leq p_m.$$

$$2^\circ \text{ If } p \leq q, \text{ then } A_{q_m} = A_q \leq A_p = A_{p_m},$$

i.e. $A_{q_m} \subseteq A_{p_m}$. Hence $q_m \vee p_m \sim q_m$ and $q_m \vee p_m \in |q_m|_{\sim}$. Thus

$$q_m \vee p_m \leq q_m, \text{ i.e. } p_m \leq q_m.$$

3° From $|p|_{\sim} = |p_m|_{\sim}$, and $(p_m)_m = \bigvee_{q \in |p_m|_{\sim}} q = \bigvee_{q \in |p|_{\sim}} q = p_m$, it follows that $(p_m)_m = p_m$.

1°, 2° and 3° prove the proposition.

Considering the set of closed elements under the closure operation defined above, we get the following corollary:

2.5. The partially ordered set $L_{\bar{A}} = \langle \{p \in L; p = p_m\}, \leq \rangle$, where the ordering relation is the one from L , is a complete lattice - the quotient relative to the closure operation $p \rightarrow p_m$.

P r o o f. This is a general property of closure operations on lattices (see, for example [5]).

To discuss the connection between \bar{A}_L and $L_{\bar{A}}$ we need the following lemma.

2.6. If $p, q \in L_{\bar{A}}$, and $p \sim q$, then $p = q$.

P r o o f. If $p \sim q$, then $p \leq q_m = q$, and $q \leq p_m = p$, i.e. $p = q$.

Let $\bar{A}_{dL} = \langle \{A_p; p \in L\}, \preceq \rangle$ be the dual lattice of \bar{A}_L . Clearly, $A_p \preceq A_q$ iff $A_q \subseteq A_p$.

2.7. $L_{\bar{A}} = \bar{A}_{dL}$.

P r o o f. Consider the mapping $f : \{p \in L; p = p_m\} \rightarrow \{A_p; p \in L\}$, such that $f(p) = A_p$.

(i) f is "onto": really, if $A_q \in \{A_p; p \in L\}$, then $A_q = A_{q_m}$, and $f(q_m) = A_{q_m} = A_q$.

(ii) f is "one to one": $f(p) = f(q)$ iff $A_p = A_q$ iff $p \sim q$ iff (by 2.6.) $p = q$.

(iii) f preserves the ordering: if $p \leq q$, then $A_q \subseteq A_p$ i.e. $A_p \preceq A_q$.

(iv) f^{-1} preserves the ordering: let $A_p \preceq A_q$, i.e. $A_q \subseteq A_p$.

Then $A_{q_m} \subseteq A_{p_m}$, since $A_p = A_{p_m}$. Thereby $A_{q_m} = A_{p_m} \cap A_{q_m} = A_{p_m} \vee q_m$ i.e. $p_m \vee q_m \in |q_m|_{\sim}$, and thus $p_m \vee q_m \leq q_m$, i.e. $p_m \leq q_m$.

By (i), (ii), (iii) and (iv), f is a lattice isomorphism.

Let now for $p \in L_{\bar{A}}$ (determined by a given fuzzy set \bar{A} on S), and for $x \in S$

$$p \cdot A_p(x) = \begin{cases} p, & \text{if } x \in A_p \\ 0 \in L_{\bar{A}}, & \text{otherwise.} \end{cases}$$

(Note that $|0|_{\sim} = 0$ (i.e. $0_{L_{\bar{A}}} = 0_L$), iff $\text{card}(|0|_{\sim}) = 1$).

2.8. If $\bar{A} = \bigcup_{p \in L} p \cdot A_p$, then $\bar{A} = \bigcup_{p \in L_{\bar{A}}} p \cdot A_p$.

P r o o f. If $x \in S$, then

$$\bar{A}(x) = \bigvee_{p \in L} p \cdot A_p(x) = \bigvee_{|p|_{\sim} \in L / \sim} \left(\bigvee_{q \in |p|_{\sim}} q \cdot A_q(x) \right) = \bigvee_{p \in L_{\bar{A}}} p \cdot A_p(x).$$

The last supremum is the same in L and in $L_{\bar{A}}$, since it is a supremum of the elements of $L_{\bar{A}}$, and the supremum itself belongs to $L_{\bar{A}}$ (otherwise, if $\bar{A}(x) = p \notin L_{\bar{A}}$, then $p < p_m$ and $A_p = A_{p_m}$. Hence $x \in A_{p_m}$, and $\bar{A}(x) \geq p_m$, which is a contradiction, proving that $\bar{A}(x) \in L_{\bar{A}}$ for every $x \in S$).

We have thus proved that for every fuzzy set $\bar{A} : S \rightarrow L$, there is a lattice $L_{\bar{A}} \subseteq L$ (being a quotient of a closure in L), such that \bar{A} is a mapping from S to $L_{\bar{A}}$, i.e. that 2.8 holds. Moreover, using the isomorphism from $L_{\bar{A}}$ to $\bar{A}_{dL} = \langle S_1, \check{\langle} \rangle$, where $S_1 = \{A_p; p \in L\} \subseteq P(S)$, one can consider \bar{A} as a mapping from S to $P(S)$.

2.9. Let $\bar{A} : S \rightarrow L$ be an arbitrary fuzzy set on S . Also let f be the isomorphism from $L_{\bar{A}}$ to \bar{A}_{dL} , defined in 2.7, and let $\check{\bar{A}} : S \rightarrow P(S)$ be the mapping for which $\check{\bar{A}}(x) \stackrel{\text{def}}{=} f(\bar{A}(x))$. Then, if $x \in S$

$$\check{\bar{A}}(x) = \bigcap_{\substack{p \in S_1 \\ x \in p}} p.$$

(Remark: If $\bar{A} = \bigcup_{p \in L_{\bar{A}}} p \cdot A_p$ and $\check{A} = \bigcup_{f(p) \in S_1} f(p) \cdot A_p^1$, then $f(\bar{A}(x)) = \check{A}(x)$).

P r o o f. $\bar{A} = \bigcup_{p \in L} p \cdot A_p = \bigcup_{p \in L_{\bar{A}}} p \cdot A_p$, by 2.8. Now, since $L_{\bar{A}} \cong \bar{A}_{dL}$, and since $f(p) = A_p$,

$$\check{A} = \bigcup_{p \in L_{\bar{A}}} f(p) \cdot A_p = \bigcup_{f(p) \in S_1} f(p) \cdot f(p).$$

Here $f(p) \cdot A_p = f(p) \cdot f(p)$ is a mapping $S \rightarrow P(S)$, such that

$$f(p) \cdot f(p)(x) = \begin{cases} f(p), & \text{if } x \in f(p) = A_p, \\ 0_{\bar{A}_{dL}} = S, & \text{otherwise,} \end{cases}$$

since $0_{\bar{A}_{dL}} = 1_{\bar{A}_L} = \bar{A}_0 = S$. Hence for $x \in S$ $A(x) = \bigvee_{\substack{p \in S \\ x \in p}} P$,

where the supremum is in \bar{A}_{dL} , the dual lattice of \bar{A}_L . Since the infimum in \bar{A}_L is the set intersection, it follows that

$$\check{A}(x) = \bigcap_{\substack{p \in S \\ x \in p}} P.$$

3. We shall now apply the results from 2 in describing the most general method of constructing the fuzzy congruence relations on an algebra A ([4]), using only the lattice of all (ordinary) congruence relations on A . This proves that, even for such a special notion as a fuzzy congruence relation, the choice of a lattice L is equivalent with the choice of a closure operation in a lattice of all congruence relation on A .

Let $A = \langle S, \theta \rangle$ be an algebra, and L an arbitrary complete lattice. The mapping $\bar{\theta} : S^2 \rightarrow L$ is a *fuzzy congruence relation* on A , iff the following conditions are satisfied ([3]):

- (r) For all $x \in S$, $\bar{\theta}(x, x) = 1 \in L$;
- (s) For all $x, y \in S$, $\bar{\theta}(x, y) = \bar{\theta}(y, x)$;
- (t) For all $x, y \in S$, $\bar{\theta}(x, y) \geq \bigvee_{z \in S} (\bar{\theta}(x, z) \wedge \bar{\theta}(z, y))$;

1) In $f(p) \cdot A_p, A_p$ is the characteristic function of A_p in the sense of 0.1.b), on the lattice \bar{A}_{dL} .

(c) If $x_1, \dots, x_n, y_1, \dots, y_n \in S$, $f \in F(n) \subseteq 0$, and

$$\bar{\theta}(x_i, y_i) = p_i, \quad i = 1, \dots, n, \quad \text{then}$$

$$\bar{\theta}(f(x_1, \dots, x_n), f(y_1, \dots, y_n)) \geq \bigwedge_{i=1}^n p_i.$$

From 2.9 we know that every fuzzy set on S can be indexed by the elements of $P(S)$, and from 1.3, if we use $P(S)$ instead of L , we get the construction of fuzzy set using $P(S)$ only. It is clear that the same method is applicable on the construction of fuzzy relations. Moreover, we claim that the properties of fuzzy congruence relations ([3]) are preserved under the described constructions.

3.1. Let $A = \langle S, 0 \rangle$ be an algebra, and let $C = \langle C, \subseteq \rangle$ be a lattice of its congruence relation. Let C_m be a subset of C closed under the arbitrary intersections, and which contains S^2 and the diagonal ϵ . Then $\bar{\theta} : S^2 \rightarrow C_m$, such that

$$\bar{\theta}(x, y) = \bigcap_{(x, y) \in \sigma} \sigma \text{ is a fuzzy congruence relation on } A.$$

Proof. $\bar{\theta}$ is a reflexive, since $\bar{\theta}(x, x) = \bigcap_{(x, x) \in \sigma} \sigma = \bigcap \tau_m = \epsilon$

(diagonal). $\bar{\theta}$ is obviously symmetric. It is transitive:

$$\bar{\theta}(x, z) \wedge \bar{\theta}(z, y) = \bigcap_{(x, z) \in \sigma_i} \sigma_i \wedge \bigcap_{(z, y) \in \tau_i} \tau_i = (\sigma \wedge \tau | (x, z) \in \sigma,$$

$(z, y) \in \tau) = (\rho | (x, z) \in \rho, (z, y) \in \rho) = (\rho | (x, y) \in \rho)$ (ρ is transitive).

$$\text{Thereby, } \bar{\theta}(x, y) = \bigcap_{(x, y) \in \rho_i} \rho_i \subseteq \rho = \bar{\theta}(x, z) \wedge \bar{\theta}(z, y).$$

Hence, $\bar{\theta}(x, y) \subseteq \bigcap_{z \in S} (\bar{\theta}(x, z) \wedge \bar{\theta}(z, y))$, i.e.

$$\bar{\theta}(x, y) \supseteq \bigvee_{z \in S} (\bar{\theta}(x, z) \wedge \bar{\theta}(z, y)).$$

$\bar{\theta}$ satisfies the substitution property (c): (we give the proof for a binary operation " \cdot ", the same is in the general case); let

$$\bar{\theta}(x_1, y_1) = \sigma_1 = \bigcap_{(x_1, y_1) \in \sigma_i} \sigma_i,$$

$$\bar{\theta}(x_2, y_2) = \sigma_2 = \bigcap_{(x_2, y_2) \in \tau_i} \tau_i. \text{ Then}$$

$\sigma_1 \wedge \sigma_2 = \tau$ and $(x_1, y_1) \in \tau$, $(x_2, y_2) \in \tau$. Since $e \in \tau_m \subseteq C$, $(x_1 \cdot x_2, y_1 \cdot y_2) \in \tau$, and hence

$$\bar{\theta}(x_1 \cdot x_2, y_1 \cdot y_2) = \bigcap_{(x_1 \cdot x_2, y_1 \cdot y_2) \in \sigma} \sigma \subseteq \tau, \text{ proving (c).}$$

Note that by 2.9 every fuzzy congruence relation on A is equal up to the lattice isomorphism, to the one constructed in 3.1, which proves that the lattice of all congruence relations on A is all we need in considering the fuzzy congruence relations on A .

REFERENCES

- [1] Kaufman A., *Introduction a la theorie des sous-ensembles flous*, Paris, 1973.
- [2] Botta O., *Theorems de representation des parties floues*, *Seminaire: Mathematique floue*, 1978-79.
- [3] Vojvodić G., Šešelja B., *O strukturi slabih relacija ekvivalencije i slabih relacija kongruencije*, *Matematički Vesnik*, 1(14)(29) 1977, 147-152.
- [4] Vojvodić G., Šešelja B., *On Fuzzy Quotient Algebras*, *Zb. rad. Prir.-mat. fak., Novi Sad, Ser. Mat.*, 13(1983), 279-288.
- [5] Algrner M., *Combinatorial theory*, Springer-Verlag, 1979.

received by the editors June 27, 1984.

IZIME

RASPLINUTI SKUPOVI NA S KAO OPERATORI ZATVARANJA NA $P(S)$

U radu se pokazuje da svaki rasplinuti skup $S \neq \emptyset$ određuje jednu kompletnu mrežu, količnik posebnog zatvorenja $P(S)$, čiji su elementi indeksirani elementima odgovarajuće

mreže L . Pokazuje se da postoji izomorfizam tako konstruisane mreže podskupova od S i količnika drugog zatvorenja - kompletne mreže sadržane u L . Tako se pokazuje da se svaki rasplinuti skup S (kao preslikavanje $S \rightarrow L$) može do na mrežni izomorfizam tretirati kao preslikavanje S u $P(S)$. Ovi rezultati primenjeni su na konstrukciju rasplnutih kongruencija na algebri A , uz pomoć mreže svih kongruencija te algebre.