

ONE CONNECTION BETWEEN BINARY  
AND  $(n+1)$ -ARY EQUIVALENCE RELATIONS ON FINITE SETS

Branimir Šešelja, Janez Ušan

Prirodno-matematički fakultet. Institut za matematiku

21000 Novi Sad, ul. dr Ilije Djuričića br.4, Jugoslavija

ABSTRACT

In this article we shall prove that every  $(n+1)$ -ary equivalence relation on a set  $S$  uniquely determines one binary equivalence in the set  $S^{\binom{n}{n}}$  of all  $n$ -subsets of  $S$ , satisfying a special property ( $(i_n)$ , Lemma 2). We shall also prove the converse, i.e. that there is a bijection between these two sets of relations. One of its consequences is (Corollary 11) that every lattice of  $(n+1)$ -ary equivalences (or of partitions of type  $n$ ) the finite set  $S$  is isomorphic to the quotient relative to one closure operation on the lattice of binary equivalences (or of partitions of type 1) on  $S^{\binom{n}{n}}$ .

Since every finite lattice is embeddable into the lattice of  $(n+1)$ -ary equivalences ( $([1])$ ), this also proves that every finite lattice is embeddable into the lattice of binary equivalences, satisfying the property  $(i_n)$  (Lemma 2).

In [1] Hartmanis defined a partition of type  $n$ , and in [2] Pickett gave a suitable definition of a corresponding  $(n+1)$ -ary equivalence relation on the same set. Some characterizations of these equivalences were given in [4]. One class of  $(n+1)$ -ary quasiordering relations and a lattice of included  $(n+1)$ -ary equivalences was considered in [5] (covering thus the contents of [3]), and the corresponding generalized orderings were discussed in [6]. Some operations on the set of all  $(n+1)$  ary relations on  $S \neq \emptyset$ , including some closure operators, were defined in [8]. In [7], it was proved that every  $(n+1)$  ary equivalence can be represented by the system of equivalences of any lower arity.

---

AMS Mathematics subject classification (1980): Primary 06A15;  
Secondary 04A05.

Key words and phrases:  $n$ -ary relations.

1. An  $(n+1)$ -ary relation  $\rho$  on the set  $S \neq \emptyset$  is  $(i, j)$ -reflexive,  $i \neq j$ ,  $i, j \in \{1, \dots, n+1\}$ , iff

$$(\forall a_1, \dots, a_{j-1}, a_{j+1}, \dots, a_{n+1} \in S) ((a_i^{i-1}, a_i, a_{i+1}^{j-1}, a_i, a_{j+1}^{n+1}) \in \rho)^1).$$

$\rho$  is reflexive iff it is  $(i, j)$ -reflexive for all  $i, j \in \{1, \dots, n+1\}$ ,  $i \neq j$  <sup>2)</sup>.

2. An  $(n+1)$ -ary relation  $\rho$  on  $S$  is  $\pi$ -symmetric,  $\pi \in \{1, \dots, n+1\}$ ! iff

$$(\forall a_1, \dots, a_{n+1} \in S) ((a_{\pi}^{n+1}) \in \rho \Rightarrow (a_{\pi(1)}, \dots, a_{\pi(n+1)}) \in \rho).$$

[2]  $\rho$  is symmetric iff it is  $\pi$ -symmetric for all  $\pi \in \{1, \dots, n+1\}$  !,

3. An  $(n+1)$ -ary relation  $\rho$  on  $S$  is  $i\bar{A}_1$ -transitive <sup>3)</sup>  $i \in \{1, \dots, n\}$ , iff

$$(\forall a_0, \dots, a_{n+1} \in S) (a_0^{i-1}, a_i, a_{i+1}^n) \in \rho \wedge (a_1^{i-1}, a_i, a_{i+1}^{n+1}) \in \rho \wedge$$

$$(a_j \neq a_k, \text{ for } j \neq k, j, k \in \{1, \dots, n\}) \Rightarrow (a_0^{i-1}, a_{i+1}^{n+1}) \in \rho).$$

4. <sub>1</sub>[1] For set  $S$  with at least  $n$  elements, the family  $\mathcal{P}_n$  of subsets of  $S$  is a partition of type  $n$ , iff

(1) each member of  $\mathcal{P}_n$  has at least  $n$  elements, and

(2) any  $n$  different elements from  $S$  belong to

exactly one number of  $\mathcal{P}_n$ .

4. <sub>2</sub>[2] An  $(n+1)$ -ary relation  $\rho$  is a generalized equivalence relation of  $S$  iff it satisfies:

E1<sub>n</sub>:  $(1, n+1)$ -reflexivity,

E2<sub>n</sub>: symmetry, and

E3<sub>n</sub>:  $n\bar{A}_1$ -transitivity.

1)  $a_p^q$  stands for  $a_p, a_{p+1}, \dots, a_{q-1}$ , and it is empty when  $q < p$ ; consequently,  $a_p^p$  is  $a_p$ , and instead of  $a, \dots, a$  ( $n$  times), we write  $\bar{a}$ ;  $\bar{a}$  is, clearly, empty.

2) In [2] an  $(n+1)$ -ary  $(1, n+1)$ -reflexive relation is called "reflexive".

3) In [2] : "transitive" stands for  $n\bar{A}_1$ -transitive.

PROPOSITION 1. Let  $S$  be a finite set,  $|S| > n$ ,  $n \in \mathbb{N}$ , and let

$$S^{(n)} = \{\{a_1, \dots, a_n\} \mid a_1, \dots, a_n \text{ are different elements from } S\}$$

If  $\rho$  is an arbitrary  $(n+1)$ -ary equivalence relation on  $S$ , let us define a binary relation  $\rho_{(n)}$  on  $S^{(n)}$ : If  $A, B \in S^{(n)}$

$$(1) \quad (A, B) \in \rho_{(n)} \quad \text{if} \quad (A \cup B)^{n+1} \subseteq \rho.$$

$\rho_{(n)}$  is a binary equivalence relation of  $S^{(n)}$ .

Proof.

1<sup>o</sup>.  $\rho_{(n)}$  is reflexive:

If  $A \in S^{(n)}$  then  $(A, A) \in \rho_{(n)}$  iff  $A^{n+1} \in \rho$ .  $|A| = n$  and thus  $(x_i^{n+1}) \in A^{n+1}$  iff  $x_i = x_j$  for some  $i \neq j$ ,  $i, j \in \{1, \dots, n+1\}$ .

$\rho$  is  $(i, j)$ -reflexive for all  $i, j \in \{1, \dots, n+1\}$  and thus  $\rho_{(n)}$  is reflexive.

2<sup>o</sup>  $\rho_{(n)}$  is symmetric:

If for  $A, B \in S^{(n)}$ ,  $(A, B) \in \rho_{(n)}$  then  $(A \cup B)^{n+1} \subseteq \rho$  i.e.  $(B \cup A)^{n+1} \subseteq \rho$ , and thus  $(B, A) \in \rho_{(n)}$ .

3<sup>o</sup>  $\rho_{(n)}$  is transitive:

Let  $A, B, C \in S^{(n)}$ ,  $(A, B) \in \rho_{(n)}$  and  $(B, C) \in \rho_{(n)}$ . Then

$$(a) \quad (A \cup B)^{n+1} \subseteq \rho \text{ and } (B \cup C)^{n+1} \subseteq \rho.$$

Consider now an arbitrary  $(n+1)$ -tuple  $(a_1^k, a_{k+1}^{n+1}) \in (A \cup C)^{n+1}$ ,

where  $a_1, \dots, a_k \in A$ ,  $c_{k+1}, \dots, c_{n+1} \in C$ ,  $k \in \{0, \dots, n+1\}$ . If

$a_1, \dots, a_k, c_{k+1}, \dots, c_{n+1}$  are not all different, then  $(a_1^k, c_{k+1}^{n+1})$

$\rho$ , since  $\rho$  is  $(i, j)$ -reflexive for all  $i, j$ . If  $a_1, \dots, a_k, c_{k+1}, \dots, c_{n+1}$  are all different, we proceed in the following way:

Let  $k = 1$ , and for all  $\alpha, \beta \in \{1, \dots, n\}$  let  $a_{\alpha(1)} \stackrel{\text{def}}{=} \bar{a}_1$ ,  $c_{\beta(1)} = \bar{c}_1, \dots, c_{\beta(n)} = \bar{c}_n$ . Now it follows from (a) that

$$(\bar{a}_1, \bar{b}_1^n) \in \rho \quad \text{and} \quad (\bar{b}_1^n, \bar{c}_1) \in \rho,$$

where  $\bar{b}_1, \dots, \bar{b}_n$  are different elements from B.

Since  $\rho$  is  $n\bar{A}_1$ -transitive and symmetric, we get  $(\bar{a}_1, \bar{b}_1^{n-1}, \bar{c}_1) \in \rho$ , that is,

$$(a1) \quad (\bar{a}_1, \bar{c}_1, \bar{b}_1^{n-1}) \in \rho.$$

From (a) we also have

$$(a2) \quad (\bar{c}_1, \bar{b}_1^{n-1}, \bar{c}_2) \in \rho.$$

The transitivity now from (a1) and (a2) gives  $(\bar{a}_1, \bar{c}_1, \bar{b}_1^{n-2}, \bar{c}_2) \in \rho$ , (we can assume here that  $\bar{c}_1 \notin \{\bar{b}_2, \dots, \bar{b}_{n-1}\}$ , since

$C \neq B$  and we choose  $\bar{c}_1$ ) and hence by the symmetry of  $\rho$

$$(a3) \quad (\bar{a}_1, \bar{c}_2, \bar{b}_1^{n-2}, \bar{c}_1) \in \rho.$$

Now we can assume that  $\bar{c}_1, \bar{c}_2, \bar{b}_1, \dots, \bar{b}_{n-2}$  are all different

(since  $\bar{c}_2$  can be equal to at least one of the  $n$  different elements from B and we omit exactly this one, say  $\bar{b}_{n-1}$ ). We now apply the transitivity on (a3) and on

$$(a4) \quad (\bar{c}_1, \bar{c}_2, \bar{b}_1^{n-2}, \bar{c}_3) \in \rho \quad (\text{by assumption (a)})$$

and we get

$$(\bar{a}_1, \bar{c}_1, \bar{c}_2, \bar{b}_1^{n-3}, \bar{c}_3) \in \rho.$$

It is now clear that this procedure in a finite number of steps gives

$$(b1) \quad (\bar{a}_1, \bar{c}_1^n) \in \rho.$$

But if we start with  $\bar{a}_2$ , we get

$$(\bar{a}_2, \bar{c}_1^n) \in \rho \quad \text{i.e.} \quad (\bar{c}_1^n, \bar{a}_2) \in \rho,$$

which, with (b<sub>1</sub>) implies  $(\bar{a}_1, \bar{c}_1^{n-1}, \bar{a}_2) \in \rho$ .

The same application of transitivity, again in a finite number of steps, finally gives

$$(a_1^{-k}, c_{k+1}^{n+1}) \in \rho,$$

proving by (1) that  $\rho_n$  is transitive.  $1^0, 2^0$  and  $3^0$  prove the proposition completely.

LEMMA 2. Let  $\rho, \rho_{(n)}$  be as in Proposition 1., and  $A, B \in S^{(n)}$ . Then,

$$(1_n) \quad (A, B) \in \rho_{(n)} \text{ iff for all } x, y \in (A \cup B)^{(n)}, (x, y) \in \rho_{(n)}.$$

Proof. If for all  $x, y \in (A \cup B)^{(n)}, (x, y) \in \rho_{(n)}$ , then it is clear that also  $(A, B) \in \rho_{(n)}$ .

If  $(A, B) \in \rho_{(n)}$ , and  $x, y \in (A \cup B)^{(n)}$ , then

$$x \cup y \subseteq A \cup B, \quad (x \cup y)^{(n)} \subseteq (A \cup B)^{(n)} \subseteq \rho, \text{ and thus by } (1_n) \\ (x, y) \in \rho_{(n)}.$$

The following proposition describes the blocks of the partition (of type 1) determined by  $\rho_{(n)}$ .

COROLLARY 3. If  $C_{\{x_1^n\}} \in S^{(n)} / \rho_{(n)}, \{x_1^n\} \in S^{(n)}$ , then

$$C_{\{x_1^n\}} = (UC_{\{x_1^n\}})^{(n) 1).$$

Proof. Let  $\{y_1^n\} \in (UC_{\{x_1^n\}})^{(n)}$ . Then,  $y_1 \in A_1, \dots, y_n \in A_n$  for some  $A_1, \dots, A_n \in C_{\{x_1^n\}}$ . From  $(A_1, A_2) \in \rho_{(n)}$ , by Lemma 2., it follows that for some  $A_1' \in (A_1 \cup A_2)^{(n)}, (A_1', A_2) \in \rho_{(n)}$ , with  $y_1, y_2 \in A_1'$ . Also, since  $(A_2, A_3) \in \rho_{(n)}$ , and since  $\rho_{(n)}$  is transitive, it follows that  $(A_1', A_3) \in \rho_{(n)}$ , and there is  $A_1'' \in (A_1' \cup A_3)^{(n)}$ , such that  $(A_1'', A_3) \in \rho_{(n)}$ , and  $y_1, y_2, y_3 \in A_1''$ . Continuing this procedure, we get  $(\{y_1^n\}, A_n) \in \rho_{(n)}$ , i.e.  $(\{y_1^n\}, \{x_1^n\}) \in \rho_{(n)}$ , and hence  $\{y_1^n\} \in C_{\{x_1^n\}}$ .

1)  $\{x_1^n\}$  is the abbreviation of  $\{x_1, \dots, x_n\}$ .

Thus,  $(\cup_{\{x_1^n\}})^{(n)} \subset C_{\{x_1^n\}}^{(n)}$ . It is obvious that  $C_{\{x_1^n\}} \subseteq (\cup_{\{x_1^n\}})^{(n)}$ ,

and thus the proof is complete.

The following two propositions prove that the connection between  $(n+1)$ -ary equivalences on  $S$  and binary equivalences on  $S^{(n)}$  is in fact a bijection.

**Proposition 4.** *If  $\rho$  is an arbitrary  $(n+1)$ -ary equivalence relation on  $S$ , and  $\rho_{(n)}$  is a binary equivalence on  $S^{(n)}$  defined in Proposition 1 (with  $(1_n)$ ), then*

$$(j_n) \quad \rho = \bigcup_{(A,B) \in \rho_{(n)}} (A \cup B)^{n+1}$$

**P r o o f.** (i) Let  $(x_1^{n+1}) \in \rho$ . If  $x_i = x_j$  for some  $i \neq j$ ,  $i, j \in \{1, \dots, n+1\}$ , then  $|\{x_1^{n+1}\}| \leq n$ , and let  $X \in S^{(n)}$ ,

$\{x_1^{n+1}\} \in X$ ,  $\rho_{(n)}$  is reflexive, and  $(X, X) \in \rho_{(n)}$ , i.e.

$$(x_1^{n+1}) \in X^{n+1} \subseteq \bigcup_{(A,B) \in \rho_{(n)}} (A \cup B)^{n+1}$$

.....  
If all  $x_1, \dots, x_{n+1}$  are different, then  $A = \{x_1^n\} \in S^{(n)}$ ,  $B = \{x_2^{n+1}\} \in S^{(n)}$  and  $(A \cup B)^{n+1} \in \rho$ . Hence, by  $(1_n)$   $(A, B) \in \rho_{(n)}$

and  $(x_1^{n+1}) \in \bigcup_{(A,B) \in \rho_{(n)}} (A \cup B)^{n+1}$ . Thus,  $\rho \subseteq \bigcup_{(A,B) \in \rho_{(n)}} (A \cup B)^{n+1}$ .

(ii) Let now  $(x_1^{n+1}) \in \bigcup_{(A,B) \in \rho_{(n)}} (A \cup B)^{n+1}$ . Then,  $(x_1^{n+1}) \in (A \cup B)^{n+1}$ , for some  $A, B \in S^{(n)}$ , such that  $(A, B) \in \rho_{(n)}$ .

Then by  $(1_n)$   $(A \cup B)^{n+1} \in \rho$ , and  $(x_1^{n+1}) \in \rho$  i.e.

$\bigcup_{(A,B) \in \rho_{(n)}} (A \cup B)^{n+1} \subseteq \rho$ . Now (i) and (ii) prove the equality  $(j_n)$ .

Proposition 5. Let  $\rho_{(n)}$  be an equivalence relation on  $S^{(n)}$ , satisfying  $(i_n)$  (Lemma 2). Then the  $(n+1)$ -ary relation  $\rho$  on  $S$ , defined by  $(j_n)$  (Corollary 4), is an  $(n+1)$ -ary equivalence relation on  $S$ .

P r o o f. (i)  $\rho$  is  $(1, n+1)$ -reflexive:

For all  $a_1, \dots, a_n \in S$ ,  $(a_1^n, a_1) \in \rho$ , since there is  $A \in S^{(n)}$  such that  $a_1, \dots, a_n \in A$ , i.e.  $(a_1^n, a_1) \in A^{n+1}$ , and  $(A, A) \in \rho_{(n)}$ .

(ii)  $\rho$  is symmetric:

If  $(a_1^{n+1}) \in \rho$  then  $(a_1^{n+1}) \in (A \cup B)^{n+1}$  for some  $A, B \in S^{(n)}$  such that  $(A, B) \in \rho_{(n)}$ , and it is clear that for every  $\pi \in \{1, \dots, n+1\}$   $(a_{\pi(1)}, \dots, a_{\pi(n+1)})$  also belongs to  $(A \cup B)^{n+1}$  i.e. to the union in  $(j_n)$ .

(iii)  $\rho$  is  $n\bar{A}_1$ -transitive:

Let  $(a_0^n) \in \rho$  and  $(a_1^{n+1}) \in \rho$ ,  $(a_1, \dots, a_n)$  are all different). Then, by  $(j_n)$  and  $(i_n)$

$$(a_0^n) \in (A \cup B)^{n+1} \quad \text{where } \{a_0^{n-1}\} \subseteq A, \{a_1^n\} = B, (A, B) \in \rho_{(n)}.$$

Then also

$$(a_1^{n+1}) \in (B \cup C)^{n+1}, \{a_2^{n+1}\} \subseteq C, (B, C) \in \rho_{(n)}.$$

But  $\rho_{(n)}$  is transitive, and thus  $(A, C) \in \rho_{(n)}$ . Hence for every  $i \in \{1, \dots, n\}$   $(a_0^{i-1}, a_{i+1}^{n+1}) \in \rho$ ,<sup>1)</sup> and of course,  $(a_0^{n-1}, a_{n+1}) \in \rho$ .

We have thus proved that there is a bijection between the set of all  $(n+1)$ -ary equivalences on  $S$  and the set of binary equivalences on  $S^{(n)}$  satisfying  $(i_n)$ .

The following corollary is its application to the corresponding partitions, from which one can derive some combinatorial results (concerning, for example, the matroids [9]).

---

1) In fact, we have proved that  $\rho$  satisfies  $1\bar{A}_1$ -transitivity.

Corollary 6. Let  $\rho$  be an arbitrary  $(n+1)$ -ary equivalence relation on  $S$ , and  $\rho_{(n)}$  a corresponding (in the sense of  $(1_n)$ ) binary equivalence on  $S^{(n)}$ . Let  $\phi_{\{x_1^n\}} \in S/\rho$  where  $y \in \phi_{\{x_1^n\}} \stackrel{\text{def}}{=} (y, x_1^n) \in \rho$ . Let also  $C_{\{x_1^n\}} \in S^{(n)}/\rho_{(n)}$

Then

$$a) \quad \phi_{\{x_1^n\}} = \cup C_{\{x_1^n\}} ;$$

$$b) \quad C_{\{x_1^n\}} = \phi_{\{x_1^n\}}^{(n)} .$$

P r o o f. a)  $y \in \phi_{\{x_1^n\}}$  iff  $(y, x_1^n) \in \rho$  iff (by  $(1_n)$ )  $(y, x_1^n) \in (A \cup B)^{n+1}$  and  $(A, B) \in \rho_{(n)}$  iff (by  $(i_n)$ )

$$y \in A \text{ and } (A, \{x_1^n\}) \in \rho_{(n)} \quad \text{iff}$$

$$y \in A \text{ and } A \in C_{\{x_1^n\}} \quad \text{iff } y \in C_{\{x_1^n\}} .$$

b) Let  $A \in S^{(n)}$ , Then  $A \in C_{\{x_1^n\}}$  iff

$$(A, \{x_1^n\}) \in \rho_{(n)} \quad \text{iff } (A \cup \{x_1^n\})^{n+1} \in \rho \quad \text{iff } A \subseteq \phi_{\{x_1^n\}} \quad \text{iff}$$

$$A \in \phi_{\{x_1^n\}}^{(n)} .$$

\* \* \*

In the following we shall construct some closure operations on the set of all binary relations on  $S^{(n)}$ , and we shall prove that if we restrict one of them to the lattice of binary equivalences, the closed elements are exactly the relations satisfying  $(i_n)$ .

For an arbitrary binary relation  $\rho_{(n)}$  on  $S^{(n)}$ ,

( $|S| \geq n$ ), define a new binary relation  $\bar{\rho}_{(n)}$  on the same set  $S^{(n)}$ :

If  $X, Y \in S^{(n)}$ ,

$(X, Y) \in \bar{\rho}_{(n)}$  iff there are  $A, B \in S^{(n)}$  such that  $X, Y \in (A \cup B)^{(n)}$

and  $(A, B) \in \rho_{(n)}$ .



LEMMA 7.  $\rho_{(n)} + \bar{\rho}_{(n)}$  is a closure operation on  $P((S^{(n)})^2)$ .

P r o o f.

$$(1) \rho_{(n)} \subseteq \bar{\rho}_{(n)}$$

Really, if  $(X, Y) \in \rho_{(n)}$  then trivially  $(X, Y) \in \bar{\rho}_{(n)}$ , (since  $X, Y \subseteq X \cup Y$ ).

$$(ii) \text{ If } \rho_{(n)} \subseteq \sigma_{(n)}, \text{ then } \bar{\rho}_{(n)} \subseteq \bar{\sigma}_{(n)}.$$

Really, let  $(X, Y) \in \bar{\rho}_{(n)}$ . Then there are  $A, B \in S^{(n)}$ , such that  $X, Y \in A \cup B$ , and  $(A, B) \in \rho_{(n)}$ . Then (since  $\rho_{(n)} \subseteq \sigma_{(n)}$ ),  $(A, B) \in \sigma_{(n)}$  and  $(X, Y) \in \bar{\sigma}_{(n)}$ .

$$(iii) \bar{\bar{\rho}}_{(n)} = \bar{\rho}_{(n)};$$

Really, if  $(X, Y) \in \bar{\bar{\rho}}_{(n)}$ , then there are  $A, B \in \bar{\rho}_{(n)}$ , such that  $X, Y \in (A \cup B)^{(n)}$  and  $(A, B) \in \bar{\rho}_{(n)}$ . But then there are  $A_1, B_1 \in S^{(n)}$ , such that  $A, B \in (A_1 \cup B_1)^{(n)}$ , and  $(A_1, B_1) \in \rho_{(n)}$ . Clearly  $X, Y \in (A_1 \cup B_1)^{(n)}$  and hence  $(X, Y) \in \bar{\rho}_{(n)}$ .

Thus  $\bar{\bar{\rho}}_{(n)} \subseteq \bar{\rho}_{(n)}$ , and with (1) it gives  $\bar{\bar{\rho}}_{(n)} = \bar{\rho}_{(n)}$ .

LEMMA 8. Let  $\rho_{(n)} \in P((S^{(n)})^2)$ . Then,  $\rho_{(n)}$  satisfies  $(i_n)$  iff  $\bar{\rho}_{(n)} = \rho_{(n)}$ .

P r o o f. If  $\rho_{(n)}$  satisfies  $(i_n)$ , then we have to prove the following logical equivalence:

$(X, Y) \in \rho_{(n)}$  iff there are  $A, B \in S^{(n)}$  such that

$$(1) \quad X, Y \in (A \cup B)^{(n)} \quad \text{and} \quad (A, B) \in \rho_{(n)}.$$

a)  $\Rightarrow$  This is trivial:  $A = X$  and  $B = Y$ ;

b)  $\Rightarrow$  Assume  $(A, B) \in \rho_{(n)}$  and  $(X, Y) \in (A \cup B)^{(n)}$ . By  $(i_n)$

$(X, Y) \in \rho_{(n)}$ .

Suppose now that  $\bar{\rho}_{(n)} = \rho_{(n)}$ . Then (1) holds. If  $(A, B) \in \rho_{(n)}$ , there are  $A_1, B_1 \in S^{(n)}$  such that  $A, B \in (A_1 \cup B_1)^{(n)}$  and  $(A_1, B_1) \in \rho_{(n)}$ . Now let  $X, Y$  be two arbitrary elements from  $(A_1 \cup B_1)^{(n)}$ . Obviously,  $X, Y \in (A_1 \cup B_1)^{(n)}$ , and thus  $(X, Y) \in \rho_{(n)}$ , i.e.  $(i_n)$  is satisfied.

LEMMA 9. If  $\rho_{(n)}$  is an equivalence relation on  $S^{(n)}$ , then

- a)  $\bar{\rho}_{(n)}$  is a tolerance relation (i.e. reflexive and symmetric) on  $S^{(n)}$ ;  
 b) If  $\sigma_{(n)}$  is an equivalence relation on  $S^{(n)}$ , also satisfying  $(i_n)$ , and if  $\rho_{(n)} \subseteq \sigma_{(n)}$ , then  $\bar{\rho}_{(n)} \subseteq \sigma_{(n)}$ .

Proof. a)  $\Delta \subseteq \rho_{(n)}$  ( $\Delta$ -diagonal on  $S^{(n)}$ ) since  $\rho_{(n)}$  is reflexive. By Lemma 7 (i)  $\rho_{(n)} \subseteq \bar{\rho}_{(n)}$  and hence  $\Delta \subseteq \bar{\rho}_{(n)}$ , i.e.  $\bar{\rho}_{(n)}$  is reflexive.

$\bar{\rho}_{(n)}$  is symmetric: if there are  $A, B \in S^{(n)}$  such that  $X, Y \in (A \cup B)^{(n)}$  and  $(A, B) \in \rho_{(n)}$ , then clearly  $(X, Y)$  and  $(Y, X)$  both belong to  $\rho_{(n)}$  (this proves that  $\bar{\rho}_{(n)}$  is symmetric, even if  $\rho_{(n)}$  is not).

b) By assumption  $\rho_{(n)} \subseteq \sigma_{(n)}$ , and hence by Lemma 7 (ii)  $\bar{\rho}_{(n)} \subseteq \bar{\sigma}_{(n)}$ . Since  $\bar{\sigma}_{(n)} = \sigma_{(n)}$  (Lemma 8), it follows that  $\bar{\rho}_{(n)} \subseteq \sigma_{(n)}$ .

Since the closure  $\bar{\rho}_{(n)}$  does not preserve the transitivity (see Example 1.), we shall define a new operation  $\rho_{(n)} \xrightarrow{\hat{\Delta}} \hat{\rho}_{(n)}$  on  $P((S^{(n)})^2)$ , as a successive application of the closure defined above ( $\bar{\rho}_{(n)}$ ) and the transitive closure of the binary relations ( $\hat{\rho}_{(n)}$ ), i.e.

$$\Delta \rho_{(n)} \stackrel{\text{def}}{=} \hat{\Delta}(\rho_{(n)})$$

Also, when applied  $i$  times, let  $\hat{\Delta}^i \rho_{(n)} = \hat{\Delta}^i \rho_{(n)}$

Proposition 10. For an arbitrary equivalence relation  $\rho_{(n)}$  on  $S^{(n)}$ , the relation

$$\hat{\rho}_{(n)} = \bigcup_{i \in \mathbb{N}} \hat{\Delta}^i \rho_{(n)}$$

is a minimal relation on  $S^{(n)}$  which satisfies  $(i_n)$  and contains  $\rho_{(n)}$  as a subrelation.

Proof. (x)  $\hat{\rho}_{(n)}$  is a relation on  $S^{(n)}$ , since for every  $i \in \mathbb{N}$   $\hat{\Delta}^i \rho_{(n)}$  is a subset of  $(S^{(n)})^2$ , and so is their union.

(xx)  $\hat{\rho}_{(n)}$  is an equivalence relation:

(r)  $\hat{\Delta} \hat{\rho}_{(n)}$ , since  $\hat{\Delta} \hat{\rho}_{(n)} \subseteq \hat{\rho}_{(n)}$ .

(s) From Lemma 9a), and from the properties of the transitive closure it follows that  $\hat{\Delta} \hat{\rho}_{(n)}$  and hence  $\hat{\rho}_{(n)}$  preserves the symmetry of  $\rho_{(n)}$ .

(t) If  $(X, Y) \in \hat{\rho}_{(n)}$ ,  $(Y, Z) \in \hat{\rho}_{(n)}$ , then for some  $i, j \in \mathbb{N}$   $(X, Y) \in \hat{\Delta}^i \rho_{(n)}$ ,  $(Y, Z) \in \hat{\Delta}^j \rho_{(n)}$ . If  $k = \max\{i, j\}$  both pairs belong to the transitive relation  $\hat{\Delta}^k \rho_{(n)}$ , and thus  $(X, Z) \in \hat{\Delta}^k \rho_{(n)}$  i.e.  $(X, Z) \in \hat{\rho}_{(n)}$ .

(xxx)  $\hat{\rho}_{(n)}$  satisfies  $(i_n)$ :

If  $(X, Y) \in \hat{\rho}_{(n)}$ , then  $(X, Y)$  belongs to some  $\hat{\Delta}^i \rho_{(n)}$  i.e. there are  $A, B \in S^{(n)}$  such that  $(A, B) \in \hat{\Delta}^i \rho_{(n)}$  and  $X, Y \in (A \cup B)^{(n)}$ . The union now has the same property.

(xxxx)  $\hat{\rho}_{(n)}$  is a minimal equivalence with the property  $(i_n)$ :

Let  $\rho_{(n)} \subseteq \sigma_{(n)}$ , where  $\sigma_{(n)}$  is an equivalence satisfying  $(i_n)$ . Then by Lemma 9,b),  $\bar{\rho}_{(n)} \subseteq \sigma_{(n)}$ , and hence  $\hat{\rho}_{(n)} \subseteq \sigma_{(n)}$  (the transitive closure). It is the same with the finite application of these operators, i.e. for every  $i \in N$ ,  $\rho_{(n)}^{\Delta^1} \subseteq \sigma_{(n)}$  implying  $\hat{\rho}_{(n)} \subseteq \sigma_{(n)}$ .

Consider now the lattice of all binary equivalences on  $S^{(n)}$ ,  $\langle E(S^{(n)}), \subseteq \rangle$ .

From the construction, it follows that  $\rho_{(n)} + \hat{\rho}_{(n)}$  is a closure operation in that lattice. The partially ordered set of the closed elements (see, for example, [9]) also a lattice (in general a complete one), so called the quotient relative to that closure operation. We denote it as follows:

$$\langle E(S^{(n)}), \subseteq \rangle \stackrel{\text{def}}{=} \{ \rho_{(n)} \mid \rho_{(n)} \subseteq E(S^{(n)}) \text{ and } \hat{\rho}_{(n)} = \rho_{(n)} \}, \subseteq \rangle .$$

(the infimum in that lattice is the same as in  $\langle E(S^{(n)}), \subseteq \rangle$ ).

Let  $\langle E_{n+1}(S), \subseteq \rangle$  be the lattice of all  $(n+1)$  ary equivalences on  $S^1$

Corollary 11.

$$\langle E_{n+1}(S), \subseteq \rangle \cong \langle \hat{E}(S^{(n)}), \subseteq \rangle .$$

**P r o o f.** From Propositions 1,4, and 5, it follows that the mapping  $f: E_{n+1}(S) \rightarrow \hat{E}(S^{(n)})$ , such that  $f(\rho) = \rho_{(n)}$  (where  $\rho_{(n)}$  is defined in  $(1_n)$ ) is 1-1 and onto. Both  $f$  and its inverse preserve the order:

Let  $\rho, \sigma \in E_{n+1}(S)$ , and  $\rho \in \sigma$ . Then

$$(a) \quad (x_1^{n+1}) \in \rho \text{ implies } (x_1^{n+1}) \in \sigma .$$

Since  $\rho = \bigcup_{(A,B) \in \rho_{(n)}} (A \cup B)^{n+1}$  and  $\sigma = \bigcup_{(A,B) \in \sigma_{(n)}} (A \cup B)^{n+1}$ ,

(a) is equivalent with

- 1) The fact that this poset is a lattice is a trivial consequence of the fact that the poset of all the partitions of type  $n$  on  $S$  is a lattice.



$$I: \rho(2), \quad II: \bar{\rho}(2), \quad III: \overset{\Delta}{\rho}(2), \quad III: \bar{\bar{\rho}}(2) = \overset{\Delta}{\bar{\rho}}(2) = \hat{\rho}(2).$$

The corresponding ternary equivalence  $\rho$  is

$$\rho = d_2 \cup \pi(a,b,c) \cup \pi(a,b,d) \cup \pi(a,c,d) \cup \pi(b,c,d) \quad 1)$$

$$S/\rho = \{\{a,b,c,d\}, \{a,e\}, \{b,e\}, \{c,e\}, \{d,e\}\}.$$

- 1)  $d_2 \stackrel{\text{def}}{=} \{(a_i^{n+1}) \mid a_1, \dots, a_{n+1} \in S \text{ and } a_i = a_j \text{ for some } i \neq j, i, j \in \{1, \dots, n+1\} \text{ (see [5])}\}.$
- $\pi(x_1^{n+1}) \stackrel{\text{def}}{=} \{(x_{\pi(1)}, \dots, x_{\pi(n+1)}) \mid \pi \in \{1, \dots, n+1\}\} \text{ (see also [5])}.$

#### REFERENCES

- [1] Hartmanis, J.: *Generalized Partitions and Lattice Embedding Theorems*, Proc. of Symposia in Pure Mathematics, Vol. II, Lattice Theory, Amer. Math. Soc. (1961) 22-30.
- [2] Pickett, H.E.: *A Note on Generalized Equivalence Relations*, Amer. Math. Monthly, 1966, 73, No. 8, 860-61.
- [3] Ušan, J., Šešelja, B., Vojvodić, G.: *Generalized Ordering and Partitions*, Matematički Vesnik, 3(16) (31), 1979, 241-47.
- [4] Ušan, J., Šešelja, B.: *Transitive n-ary Relations and Characterizations of Generalized Equivalences*, Zbornik radova Prirod.-Mat. Fak. u Novom Sadu, Ser. Mat. 11(1981), 231-245.
- [5] Šešelja, B., Ušan, J.: *Structure of Generalized Equivalences contained in  $(2, n\bar{A}1)$ -RT Relations*, Zbornik radova Prirod.-Mat. Fak. u Novom Sadu, Ser. Mat. 11(1981), 275-286.
- [6] Ušan, J., Šešelja, B.: *On Some Generalizations of Reflexive Antisymmetric, and Transitive Relations*, Proceeding of the Symposium on n-ary Structures, Skopje 1982, 175-184.
- [7] Šešelja, B., Ušan, J.: *On one Representation of Generalized Equivalences*, "Algebraic Conference", Novi Sad, 1981. 155-162.

- [8] Ušan, J., Šešelja, B.: *On Some Operations on the set  $P(S^{n+1})$* , *Prilozi MANU, Skopje, 1983, (to appear)*.
- [9] Aigner, M.: *Combinatorial Theory, Springer-Verlag, 1979.*

Received by the editors June 27, 1984.

REZIME

JEDNA VEZA IZMEDJU BINARNIH I  $(n+1)$ -ARNIH  
RELACIJA EKVIVALENCIJE NA KONAČNIM SKUPOVIMA

Uočena je veza između  $(n+1)$ -arnih ekvivalencija na konačnom skupu  $S$ , i binarnih ekvivalencija na skupu  $S^{(n)}$  svih  $n$ -podskupova od  $S$ .

Pokazano je da postoji izomorfizam između mreže  $(n+1)$ -arnih ekvivalencija i mreže količnika po posebno konstruisanom zatvorenju na mreži svih binarnih ekvivalencija na  $S^{(n)}$ .