

THE SET OF ALL WORDS OVER ALPHABET  $\{0,1\}$  OF  
LENGTH  $n$  WITH THE FORBIDDEN SUBWORD  $\underbrace{11\dots 1}_k$

*Doroslovački Rade*

*Fakultet tehničkih nauka. Institut za primenjene  
osnovne discipline, 21000 Novi Sad, V. Vlahovića 3,  
Jugoslavija*

ABSTRACT

We enumerate the number of those words of length  $n$  over the alphabet  $\{0,1\}$  in which the subword consisting of  $k$  consecutive 1's is forbidden. This number of words is counted in two different ways, which gives some new combinatorial identities.

1. DEFINITIONS AND DENOTATIONS

Let  $X$  denote a finite and nonempty set of symbols,  $X$  is called an alphabet. By  $X^n$  we shall denote the set of all strings of the length  $n$  over the alphabet  $X$ , i.e.  $X^n = \{x_1x_2\dots x_n \mid x_1, x_2, \dots, x_n \in X\}$ , the only element of  $X^0$  is the empty string, i.e. the string of length 0. The set of all finite strings over the alphabet  $X$  is  $X^* = \bigcup_{i \geq 0} X^i$ . If  $S$  is a set, then  $|S|$  is the cardinality of  $S$ . By  $\lceil n \rceil$  and  $\lfloor n \rfloor$  we denote the smallest integer  $\geq n$  and the greatest integer  $\leq n$ , respectively. By  $\ell_i(a)$  we denote the number of  $i$ 's in the string  $a \in X^*$  for  $i \in X$ ,  $\binom{n}{k} = 0$  iff  $n < k$ ,  $N_n = \{1, 2, \dots, n\}$  and

---

AMS Mathematics Subject Classification (1980): Primary 05A15.

Key words and phrases: Word, forbidden subword.

$$[x] = \begin{cases} |x| \text{ if } |x| - x \leq 0,5 \\ |x| \text{ if } |x| - x < 0,5 \end{cases}$$

## 2. RESULTS AND DISCUSSION

## THEOREM.

$$|A_k(n)| = \sum_{i_{k-1}=0}^{\lfloor \frac{(k-1)n}{k} \rfloor} \sum_{i_{k-2}=0}^{\lfloor \frac{(k-2)i_{k-1}}{k-1} \rfloor} \cdots \sum_{i_2=0}^{\lfloor \frac{2i_3}{3} \rfloor} \sum_{i_1=0}^{\lfloor \frac{i_2}{2} \rfloor} \binom{n-i_{k-1}+1}{i_{k-1}-i_{k-2}} \binom{i_{k-1}-i_{k-2}}{i_{k-2}-i_{k-3}} \cdots \binom{i_3-i_2}{i_2-i_1} \binom{i_2-i_1}{i_1}, \quad k \geq 3$$

where

$$A_k(n) = \{x \mid x = x_1 x_2 \cdots x_n \in \{0,1\}^n,$$

$$\forall i \in \mathbb{N}_{n-k+1}^{x_i x_{i+1} \cdots x_{i+k-1}} \neq \underbrace{11 \dots 1}_k\}.$$

PROOF: We shall introduce the denotation  $L_k(n) = |A_k(n)|$ . Obviously  $L_1(n) = 1$ . It is known that:

$$L_2(n) = \sum_{i_1=0}^{\lfloor n/2 \rfloor} \binom{n-i_1+1}{i_1} = \left[ \frac{1}{\sqrt{5}} \left( \frac{1+\sqrt{5}}{2} \right)^{n+2} \right]$$

(Fibonacci numbers).

We shall give the proof by induction on  $k$  ( $k \geq 3$ ) for each  $n \geq k$ . We shall prove first that the theorem is valid for  $k = 3$ , i.e.

$$L_3(n) = \sum_{i_2=0}^{\lfloor 2n/3 \rfloor} \sum_{i_1=0}^{\lfloor i_2/2 \rfloor} \binom{n-i_2+1}{i_2-i_1} \binom{i_2-i_1}{i_1}.$$

We make a partition of the set  $A_3(n)$  into subsets  $A_3^{i_2}(n)$ , where

$A_3^{i_2}(n)$  is the set of all those words of length  $n$  over the alphabet  $\{0,1\}$  which contain exactly  $i_2$  1's each, and do not contain the subword 111 i.e.

$$A_3^{i_2}(n) = \{x | x = x_1 x_2 \dots x_n \in \{0,1\}^n,$$

$$\forall i \in N_{n-2} x_i x_{i+1} x_{i+2} \neq 111, \ell_1(x) = i_2\}.$$

Let us construct the words from the set  $A_3^{i_2}(n)$ . We write one of the letters "0" or " $\lambda$ " in each of the  $i_2-1$  places between  $i_2$  1's, that is, we make some words of length  $i_2-1$  over the alphabet  $\{0,\lambda\}$ . The letter " $\lambda$ " denotes the empty letter, i.e. if the letter " $\lambda$ " is written between two 1's, then, actually, nothing is written. Since the subword 111 is forbidden in the words of the set  $A_3^{i_2}(n)$ , it follows that the set  $Q$  of the words of length  $i_2-1$  over the alphabet  $\{0,\lambda\}$  must satisfy the property that the subword  $\lambda\lambda$  is forbidden in the words of  $Q$ . Consequently

$$|Q| = |A_2(i_2-1)| = \sum_{i_1=0}^{\lfloor \frac{i_2-1}{2} \rfloor} \binom{i_2-1-i_1+1}{i_1} = \sum_{i_1=0}^{\lfloor \frac{i_2}{2} \rfloor} \binom{i_2-i_1}{i_1}$$

where  $i_1$  denotes the number of the letter  $\lambda$  in the word from the set  $Q$ . Since  $i_2-1-i_1$  0's is already written between these  $i_2$  1's, there remains to write  $n-i_2-(i_2-1-i_1) = n-2i_2+i_1+1$  0's. These 0's may be written into some of the  $i_2-1-i_1$  regions which already contain one zero each, as well as into the regions in front of and behind the word, that is into  $i_2-1-i_1+2 = i_2-i_1+1$  regions in all. We make this arrangement of  $n-2i_2+i_1+1$  0's into  $i_2-i_1+1$  regions by putting  $i_2-i_1$  compartments among these 0's. The number of permutations of these compartments and 0's equals the number of arrangements of these 0's into these regions, that is,

$$\binom{n-i_2+1}{i_2-i_1},$$

Thus

$$|A_3^{i_2}(n)| = \sum_{i_1=0}^{\lfloor \frac{i_2}{2} \rfloor} \binom{n-i_2+1}{i_2-i_1} \binom{i_2-i_1}{i_1}.$$

It is obvious from the definition of these sets  $A_3^{i_2}(n)$  that

$$A_3(n) = \bigcup_{i_2 > 0} A_3^{i_2}(n), \quad i_2 \neq i_2' \Rightarrow A_3^{i_2}(n) \cap A_3^{i_2'}(n) = \emptyset$$

and

$$i_2 > \left\lceil \frac{2n}{3} \right\rceil \Rightarrow A_3^{i_2}(n) = \emptyset.$$

Thus

$$|A_3(n)| = \sum_{i_2=0}^{\left\lceil \frac{2n}{3} \right\rceil} |A_3^{i_2}(n)| = \sum_{i_2=0}^{\left\lceil \frac{2n}{3} \right\rceil} \sum_{i_1=0}^{\lfloor \frac{i_2}{2} \rfloor} \binom{n-i_2+1}{i_2-i_1} \binom{i_2-i_1}{i_1}.$$

Let us assume that the assertion is valid for  $k$ . We shall prove, then, that it is also valid for  $k+1$ . We make a partition of the set  $A_{k+1}(n)$  into the subsets

$$A_{k+1}^{i_k}(n) = \{x | x = x_1 x_2 \dots x_n \in \{0,1\}^n,$$

$$\forall i \in N_{n-k+1} x_i x_{i+1} \dots x_{i+k-1} \neq \underbrace{11\dots 1}_{k+1}, \ell_1(x) = i_k\}.$$

The set  $A_{k+1}^{i_k}(n)$  is the set of all those words of length  $n$  over the alphabet  $\{0,1\}$ , which contain exactly  $i_k$  1's and which do not contain the subword  $\underbrace{11\dots 1}_{k+1}$ . We proceed with the construction of the set  $A_{k+1}^{i_k}(n)$ . We write one of the letters "0" or " $\lambda$ ", into each of the  $i_k - 1$  places between  $i_k$  1's, that is, we make the words of length  $i_k - 1$  over the alphabet  $\{0, \lambda\}$ . The letter  $\lambda$  denotes the empty letter, i.e., if the letter  $\lambda$  is written between two 1's, then the effect is the same as if nothing is written. Since the word  $\underbrace{11\dots 1}_{k+1}$  is forbidden in the words of the set  $A_{k+1}^{i_k}(n)$ , it follows that the subword  $\underbrace{\lambda \lambda \dots \lambda}_k$  is forbidden in the

set of words that we are making over the alphabet  $\{0,\lambda\}$ .

If we denote this set of words of length  $i_k-1$  over the alphabet  $\{0,\lambda\}$ , which do not contain the subword  $\lambda\lambda\dots\lambda$  by  $Q_1$ , then it is obvious, on the basis of the inductive hypothesis, that

$$|Q_1| = |A_k(i_k-1)| = \sum_{i_{k-1}=0}^{\binom{(k-1)(i_k-1)}{k}} \sum_{i_{k-2}=0}^{\binom{(k-2)i_{k-1}}{k-1}} \dots \sum_{i_2=0}^{\binom{2i_3}{3}} \sum_{i_1=0}^{\binom{i_2}{2}}$$

$$\binom{i_k-1-i_{k-1}+1}{i_{k-1}-i_{k-2}} \binom{i_k-1-i_{k-2}}{i_{k-2}-i_{k-3}} \dots \binom{i_3-i_2}{i_2-i_1} \binom{i_2-i_1}{i_1}$$

i.e.

$$|Q_1| = \sum_{i_{k-1}=0}^{\binom{(k-1)i_k}{k}} \sum_{i_{k-2}=0}^{\binom{(k-2)i_{k-1}}{k-1}} \dots \sum_{i_2=0}^{\binom{2i_3}{3}} \sum_{i_1=0}^{\binom{i_2}{2}}$$

$$\binom{i_k-i_{k-1}}{i_{k-1}-i_{k-2}} \binom{i_k-1-i_{k-2}}{i_{k-2}-i_{k-3}} \dots \binom{i_3-i_2}{i_2-i_1} \binom{i_2-i_1}{i_1},$$

where  $i_{k-1}$  denotes the total number of appearances of the letter  $\lambda$  in the words of the set  $Q_1$ . Since  $i_k-1-i_{k-1}$  o's is already written among these  $i_k-1$ 's, there remains still to write  $n-i_k-(i_k-1-i_{k-1}) = n-2i_k+i_{k-1}+1$  o's. These o's can be arbitrarily written into the  $i_k-1-i_{k-1}$  regions which already contain one zero each, as well as into the regions in front of and behind the word, that is, into  $i_k-1-i_{k-1}+2 = i_k-i_{k-1}+1$  regions in all. This arrangement of  $n-2i_k+i_{k-1}+1$  o's into  $i_k-i_{k-1}+1$  regions can be done in

$$\binom{n-i_k+1}{i_k-i_{k-1}} \text{ different ways.}$$

Thus

$$|A_{k+1}(n)| = \binom{n-i_k+1}{i_k-i_{k-1}} \cdot |Q_1|.$$

Since the definition of the sets  $A_{k+1}^{i_k}(n)$  and  $A_{k+1}(n)$  obviously implies

$$A_{k+1}(n) = \bigcup_{i_k > 0} A_{k+1}^{i_k}(n), \quad i_k' \neq i_k'' \Rightarrow A_{k+1}^{i_k'}(n) \cap A_{k+1}^{i_k''}(n) = \emptyset$$

and

$$i_k > \left\lceil \frac{kn}{k+1} \right\rceil \Rightarrow A_{k+1}^{i_k}(n) = \emptyset,$$

it follows

$$|A_{k+1}(n)| = \sum_{i_k=0}^{\left\lceil \frac{kn}{k+1} \right\rceil} |A_{k+1}^{i_k}(n)| = \sum_{i_k=0}^{\left\lceil \frac{kn}{k+1} \right\rceil} \binom{n-i_k+1}{i_k-i_{k-1}} |Q_1|,$$

which completes the proof.

Since

$$L_1(n) = 1 = 2^0 \cdot 1^n$$

$$L_2(n) = \left\lceil \frac{1}{\sqrt{5}} \left( \frac{1+\sqrt{5}}{2} \right)^{n+2} \right\rceil = [(1,17\dots)(1,6\dots)^n].$$

It can be counted that

$$L_3(n) = \left\lceil \frac{\alpha^{n+3}}{3\alpha^2 - 2\alpha - 1} \right\rceil = [(1,1\dots)(1,8\dots)^n]$$

where  $\alpha$  is the real root of the equation  $x^3 - x^2 - x - 1 = 0$  i.e.

$$\alpha = \frac{1}{3} \left( 1 + \sqrt[3]{19 + 3\sqrt{33}} + \sqrt[3]{19 - 3\sqrt{33}} \right).$$

Also

$$L_k(n) = C_1 \alpha_1^n + C_2 \alpha_2^n + \dots + C_k \alpha_k^n$$

where  $\alpha_i$  are the roots of the equation  $x^k - x^{k-1} - \dots - x^2 - x - 1 = 0$  for each  $i = 1, 2, \dots, k$ . The constants  $C_1, \dots, C_k$  are determined from the initial conditions.

## REFERENCES

- [1] Cvetković, D., *The generating function for variations with restrictions and paths of the graph and self-complementary graphs*, Univ. Beograd., Publ. Elektrotehnički fakultet, serija Mat. Fiz. No 320 - No 328 (1970), 27 - 34.
- [2] Einbu, J.M., *The enumeration of bit sequences that satisfy local criteria*, Publications de l'Institut Mathématique Beograd, tome 27(41) (1980) pp. 51-56.

Received by the editors November 6, 1984.

## REZIME

SKUP SVIH REČI NAD AZBUKOM  $\{0,1\}$  DUŽINE  $n$  SA  
ZABRANJENOM PODREČI  $\underbrace{11\dots 1}_k$

U ovom radu se izračunava broj svih reči dužine  $n$  nad azbukom  $\{0,1\}$ , u kojima je zabranjena podreč  $k$  uzastopnih jedinica. Ovaj broj se dobija na dva različita načina što daje neke nove kombinatorne identitete.