

ON THE PERMUTABILITY OF WEAK FUZZY  
CONGRUENCE RELATIONS

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ABSTRACT

We consider the composition "o" of fuzzy binary relations on the set  $\overline{C}_w(A)$  of weak fuzzy congruence relations on the given algebra  $A$ , using the complete lattice  $L$ . (Weak fuzzy congruence relations are defined for groupoids in [1], and it was proved in [2] that  $(\overline{C}_w(A), \leq)$  is a complete lattice, having as a homomorphic image the lattice of all fuzzy subalgebras of  $A$ ).

We prove that, provided that  $L$  is infinitely distributive,  $(\overline{C}_w(A), o)$  is a semilattice iff all weak fuzzy congruence relations on  $A$  are permutable. Permutability, on the other hand, does not imply the equality  $\overline{\rho}\overline{\theta} = \overline{\rho}\overline{V}\overline{\theta}$  in the lattice  $(\overline{C}_w(A), \leq)$ . Here we give the necessary and sufficient conditions for that equality, and we describe the connection between the operations  $o$  and  $V$  in all other cases.

1. Let  $A = (A, F)$  be an algebra,  $K \subseteq A$  the set of its constants, and  $L = (L, \wedge, \vee, 0, 1)$  a complete lattice.

a) Every mapping  $\overline{\rho} : A^2 \rightarrow L$  is a fuzzy relation on  $A$ .

b) A fuzzy relation  $\overline{\rho}$  on  $A$  is a weak fuzzy congruence relation on  $A$ , if the following is satisfied:

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If  $c \in K$ , then  $\bar{\rho}(c, c) = 1$  (weak reflexivity);

For  $x, y \in A$ ,  $\bar{\rho}(x, y) = \bar{\rho}(y, x)$  (symmetry);

For  $x, y \in A$ ,  $\bar{\rho}(x, y) > \bigvee_{z \in A} (\bar{\rho}(x, z) \wedge \bar{\rho}(z, y))$  (transitivity);

For  $x_1, \dots, x_n, y_1, \dots, y_n \in A$ ,  $f \in F_n \subseteq F$ ,

$\bar{\rho}(f(x_1, \dots, x_n), f(y_1, \dots, y_n)) > \bigwedge_{i=1}^n \bar{\rho}(x_i, y_i)$  (substitution).

If  $\overline{C_w(A)}$  is a set of all weak fuzzy congruence relations on  $A$ , then:

1.1.  $(\overline{C_w(A)}, <)$  is a complete lattice (where  $\bar{\rho} < \bar{\theta}$  iff for every  $x, y \in A$ ,  $\bar{\rho}(x, y) < \bar{\theta}(x, y)$ ). ([2]).

c) If  $\bar{\rho}$  and  $\bar{\theta}$  are two arbitrary fuzzy relations on  $A$ , then  $\bar{\rho} \circ \bar{\theta} : A^2 \rightarrow L$ , and for  $x, y \in A$ ,

$$\bar{\rho} \circ \bar{\theta}(x, y) = \bigvee_z (\bar{\rho}(x, z) \wedge \bar{\theta}(z, y)) .$$

1.2. [3] If  $L$  is complete and infinitely distributive, then  $\circ$  is an associative operation in the set of all fuzzy relations on  $A$ .

In the following,  $L$  is a complete lattice, and it is distributive if explicitly stated.

2. Here we consider the conditions for  $\circ$  to be an operation on  $\overline{C_w(A)}$ .

PROPOSITION 2.1. *Let  $L$  be an infinitely distributive lattice. Then for  $\bar{\rho}, \bar{\theta} \in \overline{C_w(A)} : \bar{\rho} \circ \bar{\theta} \in \overline{C_w(A)}$  iff  $\bar{\rho} \circ \bar{\theta} = \bar{\theta} \circ \bar{\rho}$ .*

P r o o f. Let  $\bar{\rho} \circ \bar{\theta} = \bar{\theta} \circ \bar{\rho}$ . Then, for  $c \in K$

$$\bar{\rho} \circ \bar{\theta}(c, c) = \bigvee_{z \in A} (\bar{\rho}(c, z) \wedge \bar{\theta}(z, c)) = \bar{\rho}(c, c) \wedge \bar{\theta}(c, c) = 1$$

(reflexivity)

$$\begin{aligned} \bar{\rho} \circ \bar{\theta}(x, y) &= \bigvee_z (\bar{\rho}(x, z) \wedge \bar{\theta}(z, y)) = \bigvee_z (\bar{\theta}(y, z) \wedge \bar{\rho}(z, x)) = \\ &= \bar{\theta} \circ \bar{\rho}(y, x) = \bar{\rho} \circ \bar{\theta}(y, x) \quad (\text{symmetry}). \end{aligned}$$

Composition  $\circ$  is associative, since  $L$  is infinitely distributive, and thus  $(\bar{\rho} \circ \bar{\theta}) \circ (\bar{\rho} \circ \bar{\theta}) = (\bar{\rho} \circ \bar{\rho}) \circ (\bar{\theta} \circ \bar{\theta}) < \bar{\rho} \circ \bar{\theta}$ . We use here the fact that  $\bar{\rho}_1 \subseteq \bar{\rho}$  and  $\bar{\theta}_1 \subseteq \bar{\theta}$  imply  $\bar{\rho}_1 \circ \bar{\theta}_1 \subseteq \bar{\rho} \circ \bar{\theta}$  (by the definition of  $\circ$ ). (transitivity).

Let  $f \in F$  be a binary operation denoted by " $\cdot$ ". Then, by the distributivity of  $L$ ,

$$\begin{aligned} \bar{\rho} \circ \bar{\theta}(x, y) \wedge \bar{\rho} \circ \bar{\theta}(u, v) &= (\bigvee_z (\bar{\rho}(x, z) \wedge \bar{\theta}(z, y))) \wedge (\bigvee_t (\bar{\rho}(u, t) \wedge \bar{\theta}(t, v))) \\ &= \bigvee_z \bigvee_t (\bar{\rho}(x, z) \wedge \bar{\theta}(z, y) \wedge \bar{\rho}(u, t) \wedge \bar{\theta}(t, v)) < \\ &< \bigvee_{z, t} (\bar{\rho}(x \cdot u, z \cdot t) \wedge \bar{\theta}(z \cdot t, y \cdot v)) < \bigvee_p (\bar{\rho}(x \cdot u, p) \wedge \bar{\theta}(p, y \cdot v)) = \\ &= \bar{\rho} \circ \bar{\theta}(x \cdot u, y \cdot v) . \end{aligned}$$

Hence,  $\bar{\rho} \circ \bar{\theta} \in \overline{C_w(A)}$ .

Let now  $\bar{\rho} \circ \bar{\theta} \in \overline{C_w(A)}$ . Then

$$\begin{aligned} \bar{\rho} \circ \bar{\theta}(x, y) &= \bar{\rho} \circ \bar{\theta}(y, x) = \bigvee_z (\bar{\rho}(y, z) \wedge \bar{\theta}(z, x)) = \\ &= \bigvee_z (\bar{\theta}(x, z) \wedge \bar{\rho}(z, y)) = \bar{\theta} \circ \bar{\rho}(x, y), \text{ i.e. } \bar{\rho} \circ \bar{\theta} = \bar{\theta} \circ \bar{\rho} . \quad \square \end{aligned}$$

PROPOSITION 2.2. Let  $\bar{\rho}, \bar{\theta}$  as well as  $\bar{\rho} \circ \bar{\theta} \in \overline{C_w(A)}$ . Then  $\bar{\rho} \vee \bar{\theta} = \bar{\rho} \circ \bar{\theta}$  iff for all  $x \in A$ ,  $\bar{\rho}(x, x) = \bar{\theta}(x, x)$ .

P r o o f. Let for  $x \in A$   $\bar{\rho}(x, x) = \bar{\theta}(x, x)$ . Then

$$\bar{\rho} \vee \bar{\theta} \subseteq \bar{\rho} \circ \bar{\theta} \subseteq \bar{\rho} \vee \bar{\theta} .$$

Indeed

$$\begin{aligned} \bar{\rho} \vee \bar{\theta}(x, y) &= \bar{\rho}(x, y) \vee \bar{\theta}(x, y) = (\bar{\rho}(x, y) \wedge \bar{\rho}(y, y)) \vee \\ &\vee (\bar{\theta}(x, x) \wedge \bar{\theta}(x, y)) = (\bar{\rho}(x, y) \wedge \bar{\theta}(y, y)) \vee (\bar{\rho}(x, x) \wedge \bar{\theta}(x, y)) < \\ &< \bigvee_z (\bar{\rho}(x, z) \wedge \bar{\theta}(z, y)) = \bar{\rho} \circ \bar{\theta}(x, y), \text{ since for } t, u \in A \\ &\quad \bar{\rho}(t, t) = \bar{\theta}(t, t) > \bar{\theta}(t, u) . \end{aligned}$$

Hence  $\bar{\rho} \cup \bar{\theta} \subseteq \bar{\rho} \circ \bar{\theta}$ . We have also

$\bar{\rho}(x, z) < \bar{\rho} \cup \bar{\theta}(x, z) < \bar{\tau}(x, z)$ , for every  $\bar{\tau} \in \overline{C_W(A)}$ , satisfying  $\bar{\rho} \cup \bar{\theta} \subseteq \bar{\tau}$ . Similarly,  $\bar{\theta}(z, y) < \bar{\rho} \cup \bar{\theta}(z, y) < \bar{\tau}(z, y)$ , for the same  $\bar{\tau}$ . Thus,

$$\bar{\rho} \circ \bar{\theta}(x, y) = \bigvee_z (\bar{\rho}(x, z) \wedge \bar{\theta}(z, y)) < \bigvee_z (\bar{\tau}(x, z) \wedge \bar{\tau}(z, y)) < \bar{\tau}(x, y),$$

since  $\bar{\tau}$  is transitive. Hence

$$\bar{\rho} \circ \bar{\theta} \subseteq \bigcap (\bar{\tau}; \bar{\rho} \cup \bar{\theta} \subseteq \bar{\tau}) = \bar{\rho} \vee \bar{\theta}.$$

Since  $\bar{\rho} \circ \bar{\theta} \in \overline{C_W(A)}$ ,  $\bar{\rho} \circ \bar{\theta} < \bar{\rho} \vee \bar{\theta}$  is impossible, and thus

$$\bar{\rho} \circ \bar{\theta} = \bar{\rho} \vee \bar{\theta}.$$

Let now  $\bar{\rho} \vee \bar{\theta} = \bar{\rho} \circ \bar{\theta}$ , which means that

$$\bar{\rho} \circ \bar{\theta} > \bar{\rho} \quad \text{and} \quad \bar{\rho} \circ \bar{\theta} > \bar{\theta}.$$

It is also true that

$$\bar{\rho} \circ \bar{\theta}(x, x) = \bigvee_z (\bar{\rho}(x, z) \wedge \bar{\theta}(z, x)) = \bar{\rho}(x, x) \wedge \bar{\theta}(x, x)$$

(since  $\bar{\rho}(t, t) > \bar{\rho}(t, u)$ ). Thereby

$$\bar{\rho} \circ \bar{\theta}(x, x) = \bar{\rho}(x, x) = \bar{\theta}(x, x). \quad \square$$

3. In this part we define for every subset  $T$  of  $\overline{C_W(A)}$  a special mapping  $f_T: T \rightarrow \overline{C_W(A)}$ , such that

$$\bigvee(\bar{\rho}; \bar{\rho} \in T) = \bigvee(f_T(\bar{\rho}); \bar{\rho} \in T).$$

We use  $f_T$  to describe the connection between  $\circ$  and  $\bigvee$  in  $(\overline{C_W(A)}, <)$ .

Let  $T = \{\bar{\rho}_i; i \in I\} \subseteq \overline{C_W(A)}$ ,  $I \neq \emptyset$ . For  $\bar{\rho} \in \overline{C_W(A)}$ , let

$$\bar{\delta}_{\bar{\rho}}: A^2 \rightarrow L, \quad \bar{\delta}_{\bar{\rho}}(x, y) = \begin{cases} \bar{\rho}(x, y), & \text{if } x = y \\ 0 & \text{, otherwise.} \end{cases}$$

Let also  $\bar{\Delta}_T \stackrel{\text{def}}{=} \bigvee_{i \in I} \bar{\delta}_{\bar{\rho}_i}$ , and for  $\bar{\rho} \in T$

$$(*) \quad f_T(\bar{\rho}) \stackrel{\text{def}}{=} \bigcap (\bar{\tau}; \bar{\tau} \in \overline{C_W(A)}, \bar{\Delta}_T \cup \bar{\rho} \subseteq \bar{\tau}).$$

LEMMA 3.1. If  $\bar{\rho} \in T$ , then  $\bar{\rho} < f_T(\bar{\rho})$ .

P r o o f.  $\bar{\rho} \in \bar{\Delta}_T \cup \bar{\rho} \subseteq \bar{\tau}$ , for all  $\bar{\tau}$  constituting the intersection in (\*).  $\square$

LEMMA 3.2. If  $\bar{\rho} \in T$ , then  $f_T(\bar{\rho}) < \vee_{i \in I} \bar{\rho}_i$ ;  $\bar{\rho}_i \in T$ .

P r o o f.  $\bar{\Delta}_T \cup \bar{\rho} < \vee_{i \in I} \bar{\rho}_i$ , and this supremum is one of the relations constituting the above intersection. Thus

$$f_T(\bar{\rho}) < \vee_{i \in I} \bar{\rho}_i \quad \square$$

LEMMA 3.3. a) If  $\bar{\rho} \in T$ , then  $\delta_{f_T(\bar{\rho})} = \bar{\Delta}_T$

b) If  $\bar{\rho}, \bar{\theta} \in T$ , then  $\delta_{f_T(\bar{\rho})} = \delta_{f_T(\bar{\theta})}$ .

P r o o f. a) Let  $\bar{A}_{\bar{\Delta}_T}^2 : A^2 \rightarrow L$ , such that

$$\bar{A}_{\bar{\Delta}_T}^2(x, y) = \bar{\Delta}_T(x, x) \wedge \bar{\Delta}_T(y, y).$$

Then  $\bar{A}_{\bar{\Delta}_T}^2 \in \overline{C_w(A)}$  ([2]), and  $\bar{A}_{\bar{\Delta}_T}^2(x, x) \geq \bar{\delta}_{\bar{\rho}}(x, y)$ .

Thus,  $\bar{A}_{\bar{\Delta}_T}^2$  is one of the relations constituting the intersection in (\*), and hence  $f_T(\bar{\rho}) < \bar{A}_{\bar{\Delta}_T}^2$ . That implies

$$\delta_{f_T(\bar{\rho})}(x, x) = f_T(\bar{\rho})(x, x) < \bar{A}_{\bar{\Delta}_T}^2(x, x) = \bar{\Delta}_T(x, x), \text{ i.e.}$$

$$(1) \quad \delta_{f_T(\bar{\rho})} < \bar{\Delta}_T.$$

Conversely, (\*) implies  $\bar{\Delta}_T < \bar{\tau}$ , for every  $\bar{\tau}$  from the family constituting the intersection, and

$$\bar{\Delta}_T \subseteq \cap \bar{\tau} = f_T(\bar{\rho}).$$

Thus,  $\bar{\Delta}_T(x, x) < f_T(\bar{\rho})(x, x) = \delta_{f_T(\bar{\rho})}(x, x)$ , i.e.

$$(2) \quad \bar{\Delta}_T < \delta_{f_T(\bar{\rho})}.$$

By (1) and (2)

$$\bar{\delta}_{f_T}(\bar{\rho}) = \bar{\Delta}_T .$$

b) Directly follows by a).  $\square$

PROPOSITION 3.4. If  $T = \{\bar{\rho}_i; i \in I\} \subseteq \overline{C_W(A)}$ , then

$$\bigvee_{i \in I} \bar{\rho}_i = \bigvee_{i \in I} f_T(\bar{\rho}_i) .$$

P r o o f. By Lemma 3.2., for every  $\bar{\rho} \in T$ ,  $f_T(\bar{\rho}) < \bar{\rho}$ , and

$$\bigvee_{i \in I} f_T(\bar{\rho}_i) < \bigvee_{i \in I} \bar{\rho}_i .$$

By Lemma 3.1., for every  $i \in I$ ,  $\bar{\rho}_i < f_T(\bar{\rho}_i)$ , and

$$\bigvee_{i \in I} \bar{\rho}_i < \bigvee_{i \in I} f_T(\bar{\rho}_i) .$$

Thus,  $\bigvee_{i \in I} \bar{\rho}_i = \bigvee_{i \in I} f_T(\bar{\rho}_i)$ .  $\square$

PROPOSITION 3.5. If  $L$  is an infinitely distributive lattice, and the set  $T = \{\bar{\rho}, \bar{\theta}\} \subseteq \overline{C_W(A)}$  satisfies the equality

$$f_T(\bar{\rho}) \circ f_T(\bar{\theta}) = f_T(\bar{\theta}) \circ f_T(\bar{\rho}) ,$$

then

$$\bar{\rho} \vee \bar{\theta} = f_T(\bar{\rho}) \circ f_T(\bar{\theta}) .$$

P r o o f. By Proposition 2.1.,  $f_T(\bar{\rho}) \circ f_T(\bar{\theta}) \in \overline{C_W(A)}$ , by Lemma 3.3. b)  $\bar{\delta}_{f_T}(\bar{\rho}) = \bar{\delta}_{f_T}(\bar{\theta})$ , by Proposition 2.2.

$$f_T(\bar{\rho}) \vee f_T(\bar{\theta}) = f_T(\bar{\rho}) \circ f_T(\bar{\theta}) , \text{ and by Proposition}$$

3.4. (a) is satisfied, proving the proposition.  $\square$

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## REZIME

## O PERMUTABILNOSTI SLABIH RELACIJA

## KONGRUENCIJE

U radu se ispituje kompozicija rasplinutih relacija na mreži slabih rasplinutih kongruencija date algebre. Daju se potrebni i dovoljni uslovi da to bude asocijativna operacija, kao i potrebni i dovoljni uslovi da kompozicija dve slabe rasplinite kongruencije bude supremum u odgovarajućoj mreži.